

# Zone-Based Reachability Analysis

## Lecture #17 of Advanced Model Checking

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July 3, 2009

## TCTL model checking

- Model checking timed automata against TCTL is **decidable**
  - example TCTL-formula:  $\forall \diamond^{\leq 10} goal$
- Key ingredient for decidability: finite quotient wrt. a bisimulation
  - bisimulation = equivalence on clock valuations
  - equivalence classes are called *regions*
- Region automaton is highly impractical for tool implementation
  - the number of regions lies in  $\Theta(|C|! \cdot \prod_{x \in C} c_x)$
- In practice, coarser abstractions than regions are used
  - this lecture considers time-bounded reachability using **zones**

## Reachability analysis

- **Forward** analysis:
  - starting from some initial configuration
  - determine configurations that are reachable within 1, 2, 3, . . . steps
  - until either the **goal** configuration is reached, or the computation **terminates**
- **Backward** analysis:
  - starting from the goal configuration
  - determine configurations that can reach the goal within 1, 2, 3, . . . steps
  - until either the **initial** configuration is reached, or the computation **terminates**

how can these approaches be realized for timed automata?

## Symbolic reachability analysis

- Use a **symbolic** representation of timed automata configurations

- needed as there are infinitely many configurations
- example: state regions  $\langle \ell, [\eta] \rangle$

- For set  $z$  of clock valuations and edge  $e = \ell \xrightarrow{g:\alpha,D} \ell'$  let:

$$Post_e(z) = \{ \eta' \in \mathbb{R}_{\geq 0}^n \mid \exists \eta \in z, d \in \mathbb{R}_{\geq 0}. \eta + d \models g \wedge \eta' = \text{reset } D \text{ in } (\eta + d) \}$$

$$Pre_e(z) = \{ \eta \in \mathbb{R}_{\geq 0}^n \mid \exists \eta' \in z, d \in \mathbb{R}_{\geq 0}. \eta + d \models g \wedge \eta' = \text{reset } D \text{ in } (\eta + d) \}$$

- Intuition:

- $\eta' \in Post_e(z)$  if for some  $\eta \in z$  and delay  $d$ ,  $(\ell, \eta) \xRightarrow{d} \dots \xrightarrow{e} (\ell', \eta')$
- $\eta \in Pre_e(z)$  if for some  $\eta' \in z$  and delay  $d$ ,  $(\ell, \eta) \xRightarrow{d} \dots \xrightarrow{e} (\ell', \eta')$

# Zones

- Clock constraints are *conjunctions* of constraints of the form:
  - $x \prec c$  and  $x - y \prec c$  for  $\prec \in \{ <, \leq, =, \geq, > \}$ , and  $c \in \mathbb{Z}$
- A *zone* is a set of clock valuations satisfying a clock constraint
  - a clock zone for  $g$  is the maximal set of clock valuations satisfying  $g$
- Clock zone of  $g$ :  $\llbracket g \rrbracket = \{ \eta \in \text{Eval}(C) \mid \eta \models g \}$
- The *state zone* of  $s = \langle \ell, \eta \rangle$  is  $\langle \ell, z \rangle$  with  $\eta \in z$
- For *zone*  $z$  and edge  $e$ ,  $\text{Post}_e(z)$  and  $\text{Pre}_e(z)$  are *zones*

state zones will be used as symbolic representations for configurations

# Example zones

on the black board

zones are convex polyhedra

## Operations on zones

- **Future** of  $z$ :
  - $\overrightarrow{z} = \{ \eta + d \mid \eta \in z \wedge d \in \mathbb{R}_{\geq 0} \}$
- **Past** of  $z$ :
  - $\overleftarrow{z} = \{ \eta - d \mid \eta \in z \wedge d \in \mathbb{R}_{\geq 0} \}$
- **Intersection** of two zones:
  - $z \cap z' = \{ \eta \mid \eta \in z \wedge \eta \in z' \}$
- **Clock reset** in a zone:
  - $\text{reset } D \text{ in } z = \{ \text{reset } D \text{ in } \eta \mid \eta \in z \}$
- **Inverse clock reset** of a zone:
  - $\text{reset}^{-1} D \text{ in } z = \{ \eta \mid \text{reset } D \text{ in } \eta \in z \}$

# Operations on zones: examples

on the black board

zones are closed under all aforementioned operations



## Symbolic successors and predecessors

Recall that for edge  $e = \ell \xrightarrow{g:\alpha,D} \ell'$  we have:

$$Post_e(z) = \{ \eta' \in \mathbb{R}_{\geq 0}^n \mid \exists \eta \in z, d \in \mathbb{R}_{\geq 0}. \eta + d \models g \wedge \eta' = \text{reset } D \text{ in } (\eta + d) \}$$

$$Pre_e(z) = \{ \eta \in \mathbb{R}_{\geq 0}^n \mid \exists \eta' \in z, d \in \mathbb{R}_{\geq 0}. \eta + d \models g \wedge \eta' = \text{reset } D \text{ in } (\eta + d) \}$$

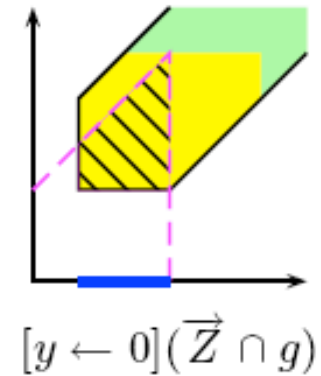
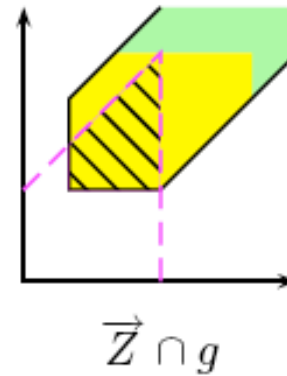
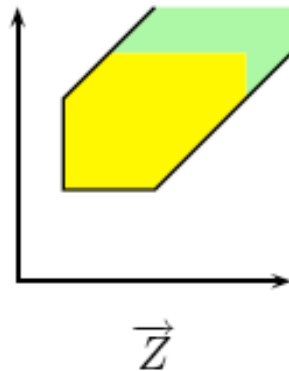
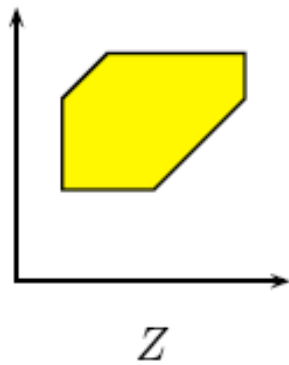
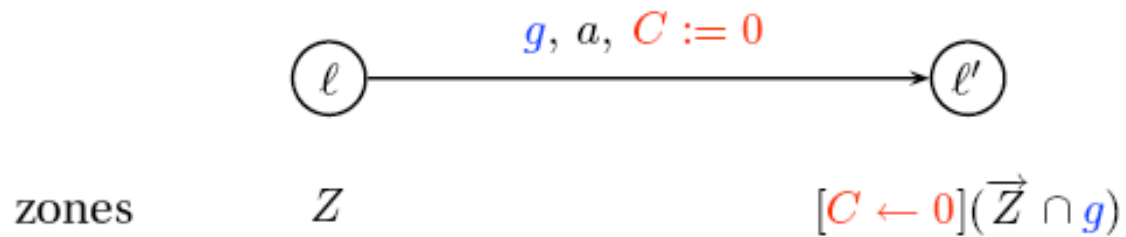
This can also be expressed symbolically using operations on zones:

$$Post_e(z) = \text{reset } D \text{ in } (\vec{z} \cap \llbracket g \rrbracket)$$

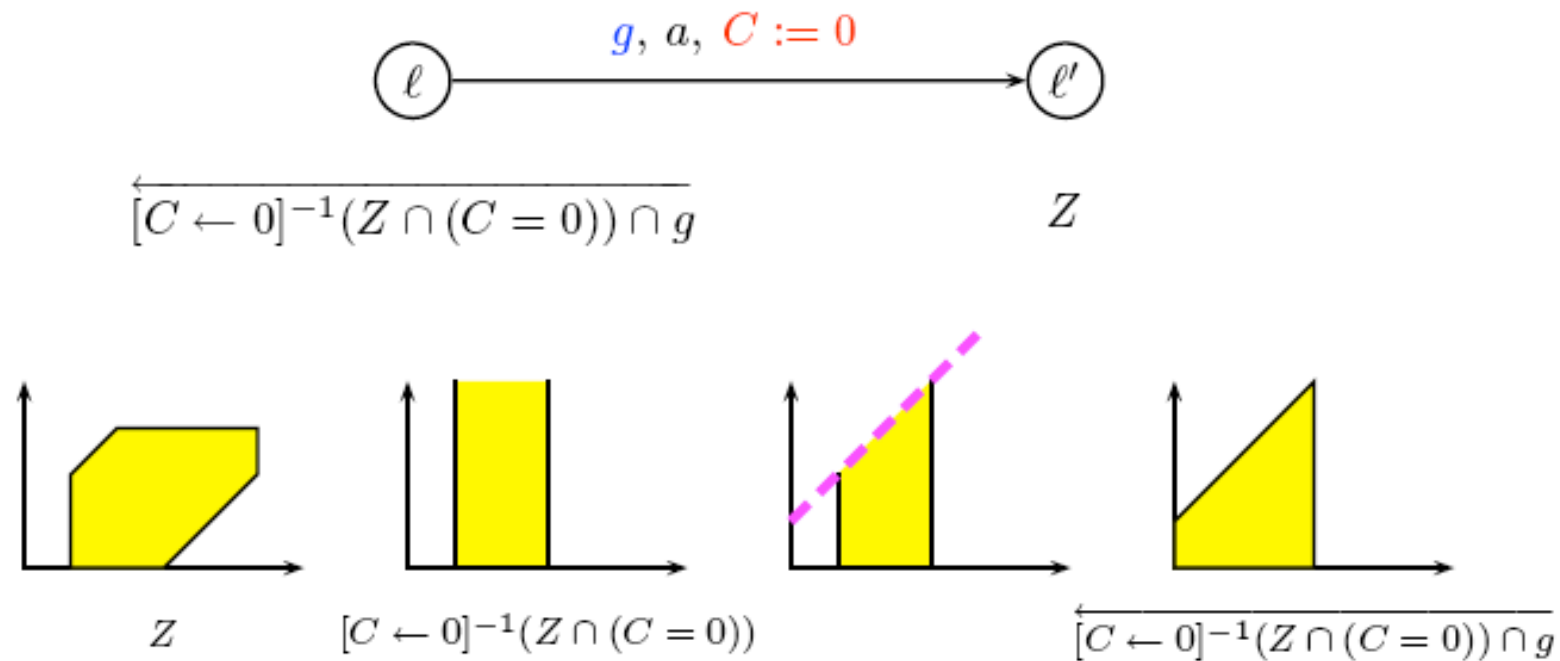
and

$$Pre_e(z) = \overleftarrow{\text{reset}^{-1} D \text{ in } (z \cap \llbracket D = 0 \rrbracket)} \cap \llbracket g \rrbracket$$

## Zone successor: example



## Zone predecessor: example



## Backward symbolic transition system (1)

Backward symbolic transition system of  $TA$  with  $|C| = n$  is inductively defined by:

$$\frac{e = \left( \ell \xrightarrow{g:\alpha,D} \ell' \right) \quad z = Pre_e(z')}{(\ell', z') \Leftarrow (\ell, z)}$$

Iterative backward reachability analysis computation schemata:

$$\begin{aligned} T_0 &= \{ (\ell, \mathbb{R}_{\geq 0}^n) \mid \ell \text{ is a goal location} \} \\ T_1 &= T_0 \cup \{ (\ell, z) \mid \exists (\ell', z') \in T_0 \text{ such that } (\ell', z') \Leftarrow (\ell, z) \} \\ \dots &\quad \dots \\ T_{k+1} &= T_k \cup \{ (\ell, z) \mid \exists (\ell', z') \in T_k \text{ such that } (\ell', z') \Leftarrow (\ell, z) \} \\ \dots &\quad \dots \end{aligned}$$

until either the computation stabilizes or reaches an initial configuration  $(\ell_0, z_0)$

## Backward symbolic transition system (2)

Backward symbolic transition system of  $TA$  is inductively defined by:

$$\frac{e = \left( \ell \xrightarrow{g:\alpha,D} \ell' \right) \quad z = Pre_e(z')}{(\ell', z') \Leftarrow (\ell, z)}$$

Iterative backward reachability analysis computation schemata:

$$\begin{aligned} T_0 &= \{ (\ell, \mathbb{R}_{\geq 0}^n) \mid \ell \text{ is a goal location} \} \\ T_1 &= T_0 \cup \{ (\ell, z) \mid \exists (\ell', z') \in T_0. (\ell', z') \Leftarrow (\ell, z) \text{ and } \ell' = \ell \text{ implies } z \not\subseteq z' \} \\ \dots &\quad \dots \\ T_{k+1} &= T_k \cup \{ (\ell, z) \mid \exists (\ell', z') \in T_k. (\ell', z') \Leftarrow (\ell, z) \text{ and } \ell' = \ell \text{ implies } z \not\subseteq z' \} \\ \dots &\quad \dots \end{aligned}$$

until either the computation stabilizes or reaches an initial configuration  $(\ell_0, z_0)$

## Termination and correctness [Henzinger et al., 1994]

The backward computation terminates and is correct wrt. reachability properties

Because of the bisimulation property, it holds:

Every set of valuations which is computed along the backward computation is a finite union of regions

## Forward reachability analysis (1)

Forward symbolic transition system of  $TA$  is inductively defined by:

$$\frac{e = \left( \ell \xrightarrow{g:\alpha,D} \ell' \right) \quad z' = \text{Post}_e(z)}{(\ell, z) \Rightarrow (\ell', z')}$$

Iterative forward reachability analysis computation schemata:

$$\begin{aligned} T_0 &= \{ (\ell_0, z_0) \mid \forall x \in C. z_0(x) = 0 \} \\ T_1 &= T_0 \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_0 \text{ such that } (\ell, z) \Rightarrow (\ell', z') \} \\ \dots &\quad \dots \\ T_{k+1} &= T_k \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_k \text{ such that } (\ell, z) \Rightarrow (\ell', z') \} \\ \dots &\quad \dots \end{aligned}$$

until either the computation stabilizes or reaches a symbolic state containing a goal configuration

## Forward reachability analysis (2)

Forward symbolic transition system of  $TA$  is inductively defined by:

$$\frac{e = \left( \ell \xrightarrow{g:\alpha,D} \ell' \right) \quad z' = Post_e(z)}{(\ell, z) \Rightarrow (\ell', z')}$$

Iterative forward reachability analysis computation schemata:

$$T_0 = \{ (\ell_0, z_0) \mid \forall x \in C. z_0(x) = 0 \}$$

$$T_1 = T_0 \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_0. (\ell, z) \Rightarrow (\ell', z') \text{ and } \ell = \ell' \text{ implies } z \not\subseteq z' \}$$

...

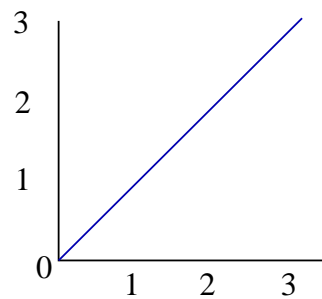
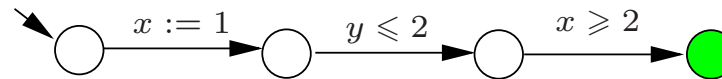
$$T_{k+1} = T_k \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_k. (\ell, z) \Rightarrow (\ell', z') \text{ and } \ell = \ell' \text{ implies } z \not\subseteq z' \}$$

...

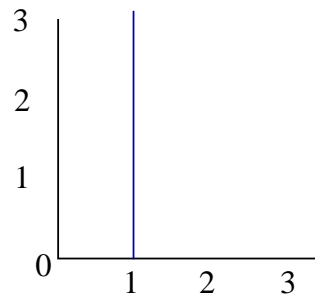
until either the computation stabilizes or reaches a symbolic state containing a goal configuration



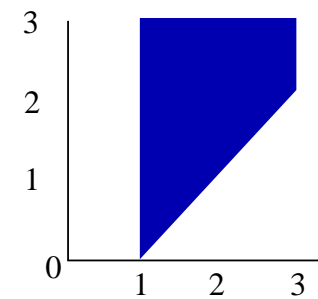
# Forward reachability analysis: intuition



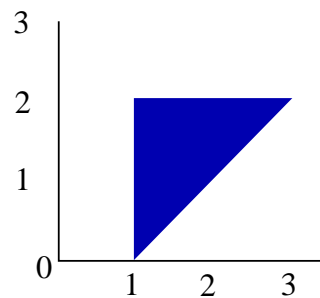
leaving initial



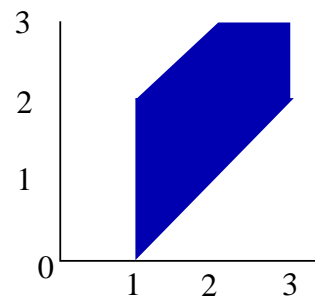
entering first



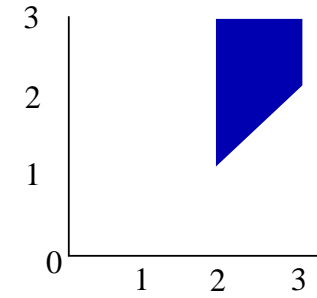
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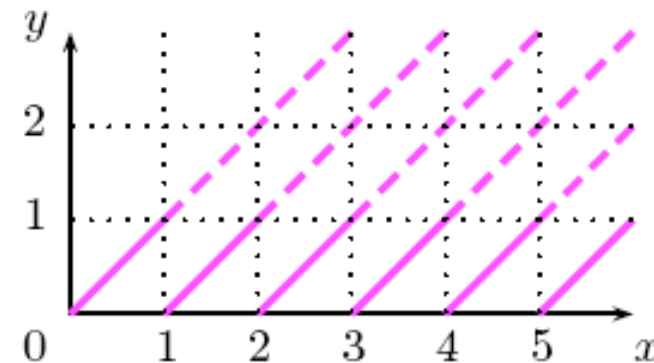
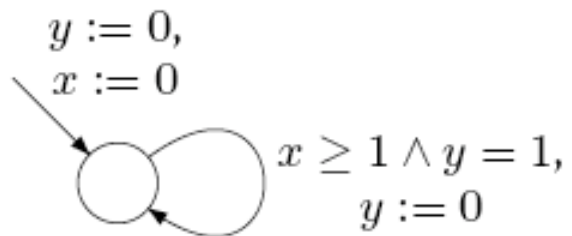
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entering third

## Possible non-termination

The forward analysis is correct but may **not** terminate:



➔ an infinite number of steps...

## Solution: abstract forward reachability

Let  $\gamma$  associate sets of valuations to sets of valuations

**Abstract** forward symbolic transition system of  $TA$  is defined by:

$$\frac{(\ell, z) \Rightarrow (\ell', z') \quad z = \gamma(z)}{(\ell, z) \Rightarrow_{\gamma} (\ell', \gamma(z'))}$$

Iterative forward reachability analysis computation schemata:

$$\begin{aligned} T_0 &= \{ (\ell_0, \gamma(z_0)) \mid \forall x \in C. z_0(x) = 0 \} \\ T_1 &= T_0 \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_0 \text{ such that } (\ell, z) \Rightarrow_{\gamma} (\ell', z') \} \\ \dots &\quad \dots \\ T_{k+1} &= T_k \cup \{ (\ell', z') \mid \exists (\ell, z) \in T_k \text{ such that } (\ell, z) \Rightarrow_{\gamma} (\ell', z') \} \\ \dots &\quad \dots \end{aligned}$$

with inclusion check and termination criteria as before