

Compiler Construction

Lecture 1: Introduction

Thomas Noll

Lehrstuhl für Informatik 2
(Software Modeling and Verification)

RWTH Aachen University

`noll@cs.rwth-aachen.de`

`http://www-i2.informatik.rwth-aachen.de/i2/cc10/`

Winter semester 2010/11

1 Preliminaries

2 Introduction

- Lectures: **Thomas Noll**
 - Lehrstuhl für Informatik 2, Room 4211
 - E-mail `noll@cs.rwth-aachen.de`
 - Phone (0241)80-21213
- Exercise classes: **Christina Jansen**
 - E-mail `christina.jansen@cs.rwth-aachen.de`
- Student assistants:
 - **Stefan Breuer**
 - **Ernst Wrtal**

- **BSc Informatik:** V3 Ü2, 6 credits
 - Wahlpflichtfach Theorie
- **MSc Informatik:** V3 Ü2, 6 credits
 - Theoretische Informatik
- **MSc Software Systems Engineering:** V4 Ü2, 8 credits
 - Theoretical CS
 - Specialization *Formal Methods, Programming Languages and Software Validation*
- **Diplomstudiengang Informatik:** V4 Ü2
 - Theoretische (+ Praktische) Informatik
 - Vertiefungsfach *Formale Methoden, Programmiersprachen und Softwarevalidierung*
 - Combination with Katoen, Thomas, Vöcking, ...; Kobbelt, Seidl, ...

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 - application of theoretical concepts
 - compiler = example of a complex software architecture
 - gaining experience with tool support

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- What **we** expect: basic knowledge in
 - imperative programming languages
 - formal languages and automata theory

- **Schedule:**

- Lecture Tue 14:00–15:30 AH 2 (starting October 19)
- Lecture Thu 13:30–15:00 AH 1 (starting October 14)
- Exercise class Wed 10:00–11:30 AH 2 (starting October 20)
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- Written material in **English**, lecture and exercise classes in **German**, rest up to you

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What Is It All About?

Compiler = Program: Source code $\rightarrow$ Target code

Source code: in **high-level programming language**, tailored to problem

- imperative vs. declarative (functional, logic) vs. object-oriented
- sequential vs. concurrent

Target code: usually **machine code**

- architecture dependent (RISC/CISC/parallel)

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More applications of compiler techniques:

- Parsing of structured data (HTML, XML, ...)
- Cross-compiling: Java $\rightarrow$ C
- File conversion: $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X} \rightarrow$ PDF
- PostScript interpreters
- ...

Properties of a Good Compiler

Correctness

Goals: **conformance** to source and target language specifications;

“equivalence” of source and target code

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Remark: **mutual tradeoffs!**

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Pragmatics

- length and understandability of programs
- learnability of programming language
- appropriateness for specific applications
- ...

Example

- ① From NASA's Mercury Project: FORTRAN DO loop
 - `DO 5 K = 1,3`: DO loop with index variable `K`
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- 3 What if `p = nil` in the following program?

```
while p <> nil and p^.key < val do ...
```

Pascal: strict Boolean operations



Modula: non-strict Boolean operations



Microsoft Visual C++ Debug Library



Debug Error!

Program: c:\code\oreilly\security\security\debug\Security.exe

A buffer overrun has been detected which has corrupted the program's internal state. The program cannot safely continue execution and must now be terminated.

(Press Retry to debug the application)

Abort

Retry

Ignore

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Code generation: since 1940s

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Automatic compiler generation: since 1980s

- [f]lex, yacc, ANTLR, action semantics, ...
- see <http://catalog.compilertools.net/>

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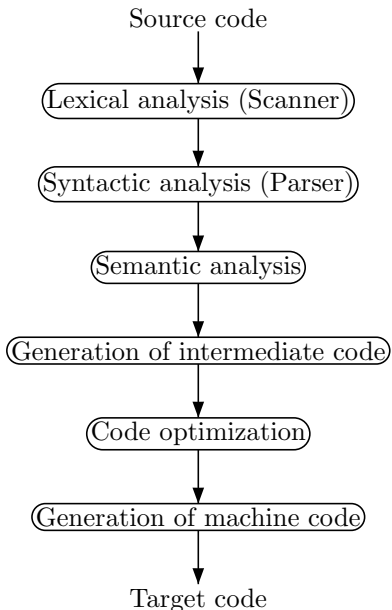
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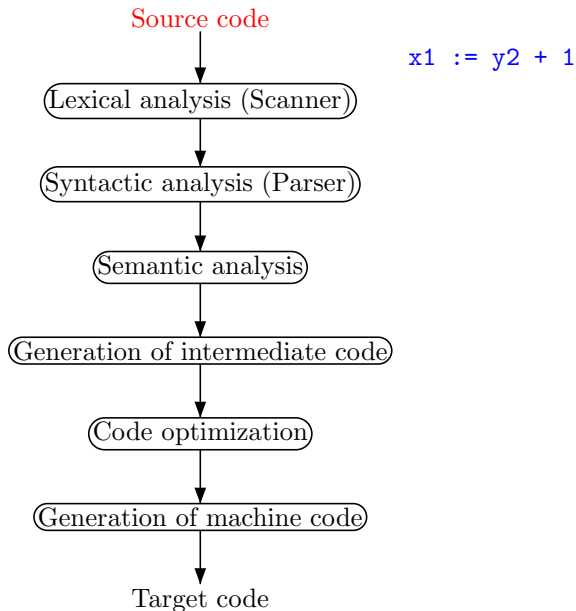
Generation of target code: tailored to target system

Additionally: optimization of target code, symbol table, error handling

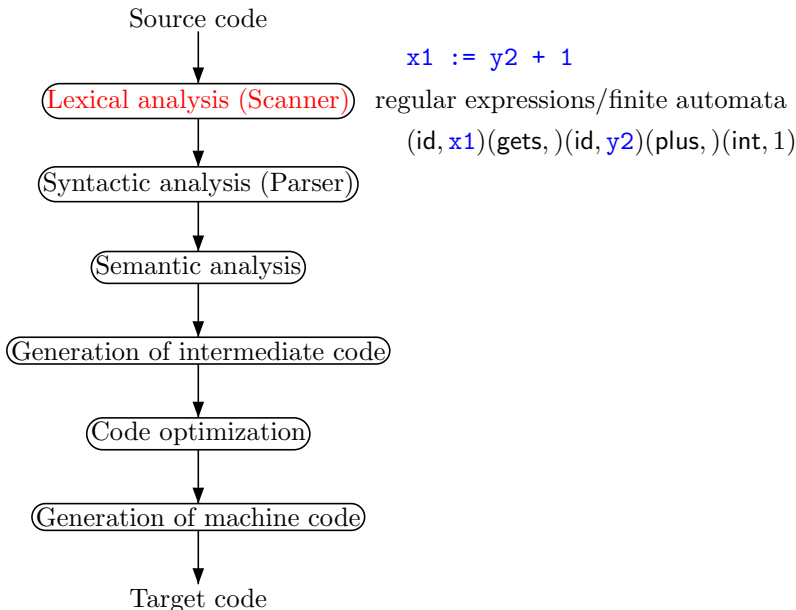
Conceptual Structure of a Compiler



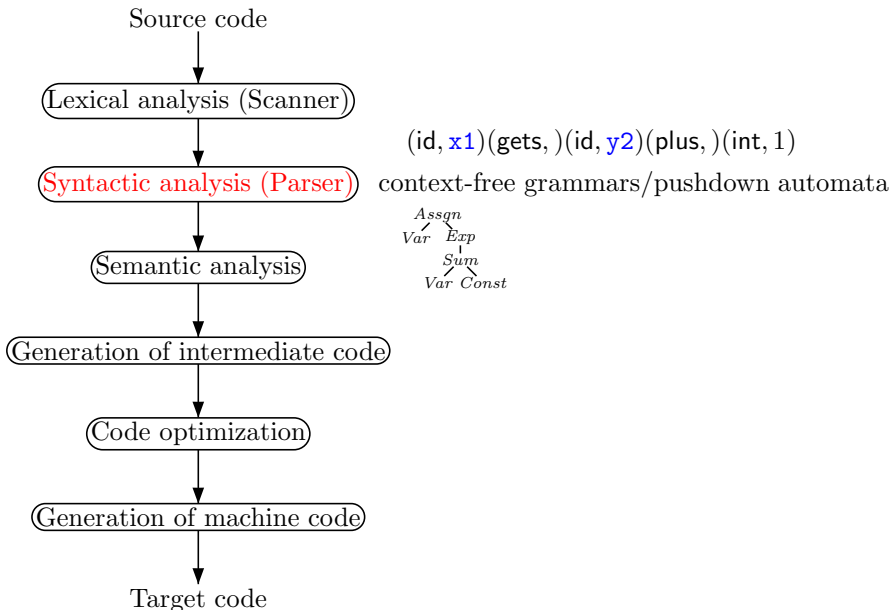
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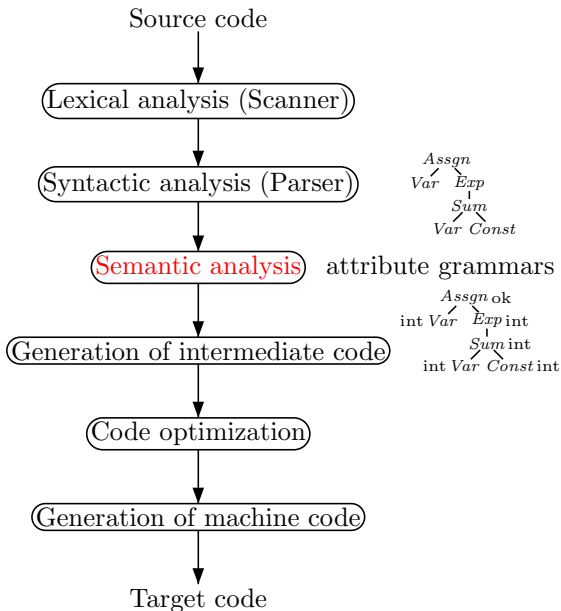
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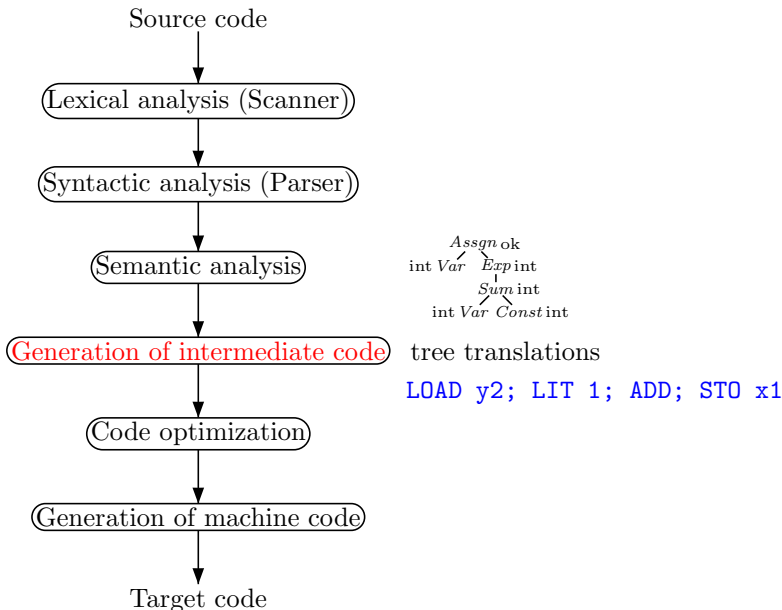
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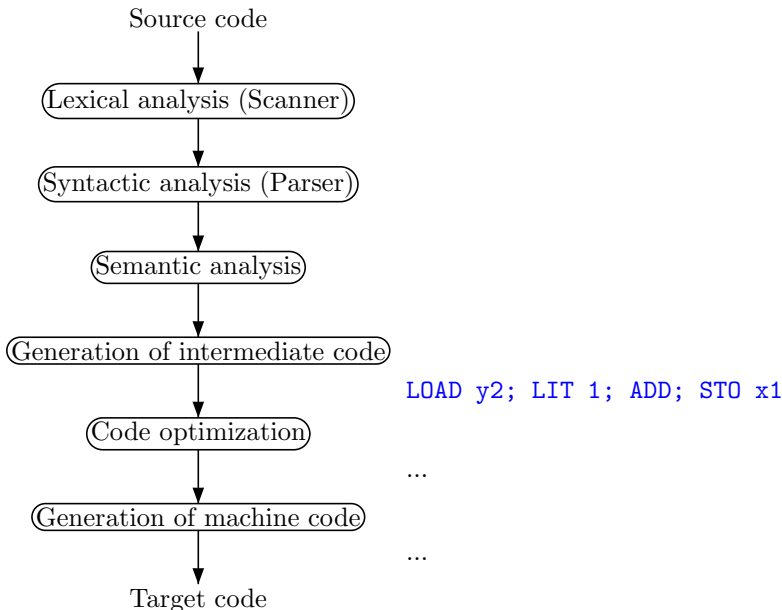
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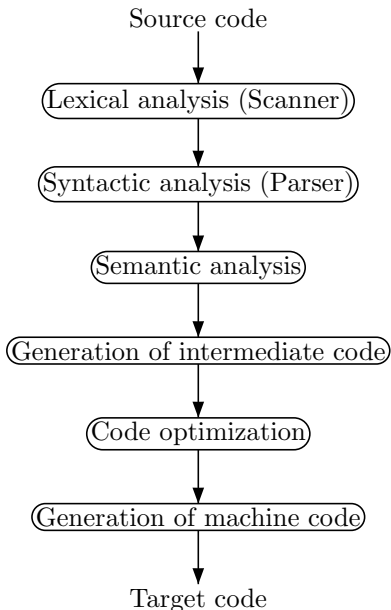
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[omitted: symbol table, error handling]

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Synthesis: generation of (intermediate/machine) code + optimization

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Historical: *n*-pass compiler

- n = number of runs through source program
- nowadays mainly one-pass

(CS Library: “Handapparat *Programmiersprachen und Verifikation*”)

General

- A.V. Aho, M.S. Lam, R. Sethi, J.D. Ullman: *Compilers – Principles, Techniques, and Tools*; 2nd ed., Addison-Wesley, 2007
- A.W. Appel, J. Palsberg: *Modern Compiler Implementation in Java*, Cambridge University Press, 2002
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