

Compiler Construction

Lecture 1: Introduction

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(Software Modeling and Verification)

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`http://www-i2.informatik.rwth-aachen.de/i2/cc10/`

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1 Preliminaries

2 Introduction

- Lectures: **Thomas Noll**
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- Exercise classes: **Christina Jansen**
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 - **Stefan Breuer**
 - **Ernst Wrtal**

- **BSc Informatik:** V3 Ü2, 6 credits
 - Wahlpflichtfach Theorie
- **MSc Informatik:** V3 Ü2, 6 credits
 - Theoretische Informatik
- **MSc Software Systems Engineering:** V4 Ü2, 8 credits
 - Theoretical CS
 - Specialization *Formal Methods, Programming Languages and Software Validation*
- **Diplomstudiengang Informatik:** V4 Ü2
 - Theoretische (+ Praktische) Informatik
 - Vertiefungsfach *Formale Methoden, Programmiersprachen und Softwarevalidierung*
 - Combination with Katoen, Thomas, Vöcking, ...; Kobbelt, Seidl, ...

- What **you** can expect:
 - how to implement (imperative) programming languages
 - application of theoretical concepts
 - compiler = example of a complex software architecture
 - gaining experience with tool support
- What **we** expect: basic knowledge in
 - imperative programming languages
 - formal languages and automata theory

- **Schedule:**
 - Lecture Tue 14:00–15:30 AH 2 (starting October 19)
 - Lecture Thu 13:30–15:00 AH 1 (starting October 14)
 - Exercise class Wed 10:00–11:30 AH 2 (starting October 20)
 - see overview at <http://www-i2.informatik.rwth-aachen.de/i2/cc10/>
- **1st assignment sheet** next week, presented October 27
- Work on assignments in **groups of three**
- **Written exam** on Tue February 1
 - for BSc/MSc candidates (6/8 credits)
 - for Diplom candidates (Übungsschein)
- **Admission** requires at least 50% of the points in the exercises
- Written material in **English**, lecture and exercise classes in **German**, rest up to you

1 Preliminaries

2 Introduction

What Is It All About?

Compiler = Program: Source code \rightarrow Target code

Source code: in **high-level programming language**, tailored to problem

- imperative vs. declarative (functional, logic) vs. object-oriented
- sequential vs. concurrent

Target code: usually **machine code**

- architecture dependent (RISC/CISC/parallel)

More applications of compiler techniques:

- Parsing of structured data (HTML, XML, ...)
- Cross-compiling: Java \rightarrow C
- File conversion: $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X} \rightarrow \text{PDF}$
- PostScript interpreters
- ...

Properties of a Good Compiler

Correctness

Goals: **conformance** to source and target language specifications;
“**equivalence**” of source and target code

- compiler validation and verification
- proof-carrying code, ...

Efficiency of generated code

Goal: target code as **fast** and/or **memory efficient** as possible

- program analysis and optimization

Efficiency of compiler

Goal: translation process as **fast** and/or **memory efficient** as possible
(for inputs of arbitrary size)

- fast (linear-time) algorithms
- sophisticated data structures

Remark: **mutual tradeoffs!**

Aspects of a Programming Language

Syntax: “How does a program look like?”

- hierarchical composition of programs from structural components

Semantics: “What does this program mean?”

“Static semantics”: properties which are not (easily) definable in syntax

(declaredness of identifiers, type correctness, ...)

“Dynamic semantics”: execution evokes state transformations of an (abstract) machine

Pragmatics

- length and understandability of programs
- learnability of programming language
- appropriateness for specific applications
- ...

Example

- 1 From NASA's Mercury Project: FORTRAN DO loop
 - `DO 5 K = 1,3`: DO loop with index variable `K`
 - `DO 5 K = 1.3`: assignment to (`real`) variable `D05K`
- 2 How often is the following loop traversed?

```
for i := 2 to 1 do ...
```

FORTRAN IV: once

PASCAL: never

- 3 What if `p = nil` in the following program?

```
while p <> nil and p^.key < val do ...
```

Pascal: strict Boolean operations



Modula: non-strict Boolean operations



Microsoft Visual C++ Debug Library



Debug Error!

Program: c:\code\oreilly\security\security\debug\Security.exe

A buffer overrun has been detected which has corrupted the program's internal state. The program cannot safely continue execution and must now be terminated.

(Press Retry to debug the application)

Abort

Retry

Ignore

Historical Development

Code generation: since 1940s

- ad-hoc techniques
- concentration on back-end
- first FORTRAN compiler in 1960

Formal syntax: since 1960s

- LL/LR parsing
- shift towards front-end
- semantics defined by compiler/interpreter

Formal semantics: since 1970s

- operational
- denotational
- axiomatic
- see course *Semantics and Verification of Software*

Automatic compiler generation: since 1980s

- [f]lex, yacc, ANTLR, action semantics, ...
- see <http://catalog.compilertools.net/>

Lexical analysis (Scanner):

- recognition of symbols, delimiters, and comments
- by regular expressions and finite automata

Syntactic analysis (Parser):

- determination of hierarchical program structure
- by context-free grammars and pushdown automata

Semantic analysis:

- checking context dependencies, data types, ...
- by attribute grammars

Generation of intermediate code:

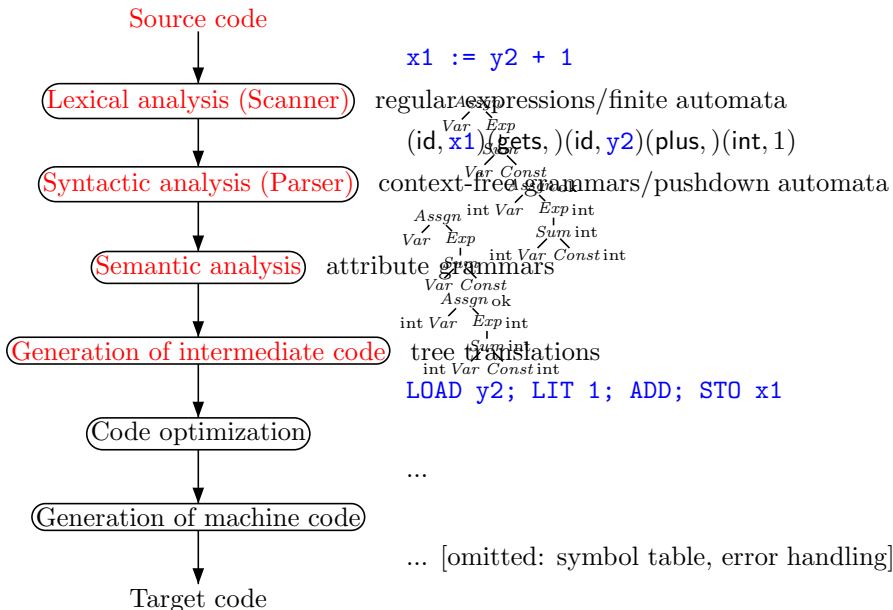
- translation into (target-independent) intermediate code
- by tree translations

Code optimization: to improve runtime and/or memory behavior

Generation of target code: tailored to target system

Additionally: optimization of target code, symbol table, error handling

Conceptual Structure of a Compiler



Classification of Compiler Phases

Analysis vs. synthesis

Analysis: lexical/syntactic/semantic analysis
(determination of syntactic structure, error handling)

Synthesis: generation of (intermediate/machine) code + optimization

Front-end vs. back-end

Front-end: machine-independent parts
(analysis + intermediate code + machine-independent optimizations)

Back-end: machine-dependent parts
(generation + optimization of machine code)

Historical: *n*-pass compiler

- n = number of runs through source program
- nowadays mainly one-pass

(CS Library: “Handapparat Programmiersprachen und Verifikation”)

General

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Special

- O. Mayer: *Syntaxanalyse*, BI-Wissenschafts-Verlag, 1978
- D. Brown, R. Levine T. Mason: *lex & yacc*, O'Reilly, 1995
- T. Parr: *The Definite ANTLR Reference*, Pragmatic Bookshelf, 2007

Historical

- W. Waite, G. Goos: *Compiler Construction, 2nd edition*, Springer, 1985
- N. Wirth: *Grundlagen und Techniken des Compilerbaus*, Addison-Wesley, 1996