

Compiler Construction

Lecture 20: Code Generation III (Translation to Intermediate Code)

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- 1 Repetition: Intermediate Code
- 2 The Symbol Table
- 3 Translation of Programs
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Definition (Syntax of EPL)

The **syntax of EPL** is defined as follows:

$\mathbb{Z} :$ z (* z is an integer *)

$Ide :$ I (* I is an identifier *)

$AExp :$ $A ::= z \mid I \mid A_1 + A_2 \mid \dots$

$BExp :$ $B ::= A_1 < A_2 \mid \text{not } B \mid B_1 \text{ and } B_2 \mid B_1 \text{ or } B_2$

$Cmd :$ $C ::= I := A \mid C_1; C_2 \mid \text{if } B \text{ then } C_1 \text{ else } C_2 \mid$
 $\text{while } B \text{ do } C \mid I()$

$Dcl :$ $D ::= D_C D_V D_P$
 $D_C ::= \varepsilon \mid \text{const } I_1 := z_1, \dots, I_n := z_n;$
 $D_V ::= \varepsilon \mid \text{var } I_1, \dots, I_n;$
 $D_P ::= \varepsilon \mid \text{proc } I_1; K_1; \dots; I_n; K_n;$

$Blk :$ $K ::= D C$

$Pgm :$ $P ::= \text{in/out } I_1, \dots, I_n; K.$

The Abstract Machine AM

Definition (Abstract machine for EPL)

The **abstract machine for EPL (AM)** is defined by the **state space**

$$S := PC \times DS \times PS$$

with

- the **program counter** $PC := \mathbb{N}$,
- the **data stack** $DS := \mathbb{Z}^*$ (top of stack to the right), and
- the **procedure stack** (or: **runtime stack**) $PS := \mathbb{Z}^*$ (top of stack to the left).

Thus a state $s = (l, d, p) \in S$ is given by

- a program label $l \in PC$,
- a data stack $d = d.r : \dots : d.1 \in DS$, and
- a procedure stack $p = p.1 : \dots : p.t \in PS$.

Definition (AM instructions)

The set of **AM instructions** is divided into

arithmetic instructions: **ADD**, **MULT**, ...

Boolean instructions: **NOT**, **AND**, **OR**, **LT**, ...

jumping instructions: **JMP**(*ca*), **JFALSE**(*ca*) ($ca \in PC$)

procedure instructions: **CALL**(*ca*, *dif*, *loc*) ($ca \in PC$, $dif, loc \in \mathbb{N}$), **RET**

transfer instructions: **LOAD**(*dif*, *off*), **STORE**(*dif*, *off*) ($dif, off \in \mathbb{N}$),
LIT(*z*) ($z \in \mathbb{Z}$)

Structure of Procedure Stack I

The semantics of procedure and transfer instructions requires a particular structure of the procedure stack $p \in PS$: it must be composed of **frames** (or: **activation records**) of the form

$$sl : dl : ra : v_1 : \dots : v_k$$

where

static link sl : points to frame of surrounding declaration environment
 $\implies$ used to access non-local variables

dynamic link dl : points to previous frame (i.e., of calling procedure)
 $\implies$ used to remove topmost frame after termination of procedure call

return address ra : program label after termination of procedure call
 $\implies$ used to continue program execution after termination of procedure call

local variables v_i : values of locally declared variables

Structure of Procedure Stack II

Observation:

- The usage of a variable in a procedure body refers to its **innermost declaration**.
- If the level difference between the usage and the declaration is *dif*, then a **chain of *dif* static links** has to be followed to access the corresponding frame.

Structure of Procedure Stack II

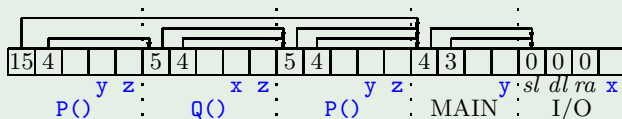
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Example (cf. Example 19.5)

```
in/out x;  
const c = 10;  
var y;  
proc P;  
  var y, z;  
  proc Q;  
    var x, z;  
    [... P() ...]  
  [... x ... y ... Q() ...]  
  proc R;  
    [... P() ...]  
  [... P() ...].
```

Procedure stack after second call of P:



Structure of Procedure Stack II

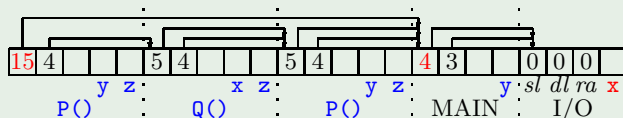
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  [... x ... y ... Q() ...]  
  proc R;  
    [... P() ...]  
  [... P() ...].
```

Procedure stack after second call of P:



P uses x $\Rightarrow$ *dif* = 2

Structure of Procedure Stack II

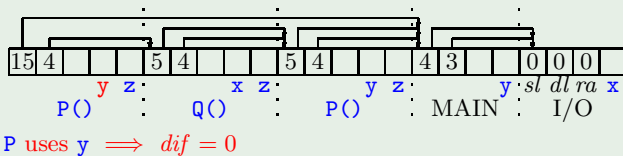
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proc R;  
  [... P() ...]  
  [... P() ...].
```

Procedure stack after second call of P:



Definition (Semantics of AM programs)

An **AM program** is a sequence of $k \geq 1$ labeled AM instructions:

$$P = 1 : O_1; \dots; k : O_k$$

The set of all AM programs is denoted by AM .

The **semantics of AM programs** is determined by

$$\llbracket . \rrbracket : AM \times S \dashrightarrow S$$

with

$$\llbracket P \rrbracket(l, d, p) := \begin{cases} \llbracket P \rrbracket(\llbracket O_l \rrbracket(l, d, p)) & \text{if } l \in [k] \\ (l, d, p) & \text{otherwise} \end{cases}$$

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Translation of EPL into AM Programs

Goal: define translation mapping

$$\text{trans} : Pgm \dashrightarrow AM$$

Translation of EPL into AM Programs

Goal: define **translation mapping**

$$\text{trans} : Pgm \dashrightarrow AM$$

The translation employs a **symbol table**:

$$\begin{aligned} Tab := \{st \mid st : Ide \dashrightarrow (\{ \text{const} \} \times \mathbb{Z}) \\ \cup (\{ \text{var} \} \times Lev \times Off) \\ \cup (\{ \text{proc} \} \times PC \times Lev \times Size)) \} \end{aligned}$$

whose entries are created by declarations:

- constant declarations: (const, z)
 - **value** $z \in \mathbb{Z}$
- variable declarations: (var, lev, off)
 - **declaration level** $lev \in Lev := \mathbb{N}$ ($0 \cong \text{I/O}$, $1 \cong \text{MAIN}$, ...)
 - **offset** $off \in Off := \mathbb{N}$
 - offset and difference between usage and declaration level determine procedure stack entry
- procedure declarations: $(\text{proc}, ca, lev, loc)$
 - **code address** $ca \in PC$
 - **declaration level** $lev \in Lev$
 - **number of local variables** $loc \in Size := \mathbb{N}$

Maintaining the Symbol Table

The symbol table is maintained by the function $\text{update}(D, \text{st}, l)$ which specifies the update of symbol table st according to declaration D (with respect to current level l):

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Definition 20.1 (update function)

$$\text{update} : Dcl \times Tab \times Lev \dashrightarrow Tab$$

is defined by

$$\begin{aligned} & \text{update}(D_C \ D_V \ D_P, \text{st}, l) \\ & \quad := \text{update}(D_P, \text{update}(D_V, \text{update}(D_C, \text{st}, l), l), l) \\ & \quad \quad \text{if all identifiers in } D_C \ D_V \ D_P \text{ different} \\ & \text{update}(\varepsilon, \text{st}, l) \\ & \quad := \text{st} \\ & \text{update}(\text{const } I_1 := z_1, \dots, I_n := z_n; \text{st}, l) \\ & \quad := \text{st}[I_1 \mapsto (\text{const}, z_1), \dots, I_n \mapsto (\text{const}, z_n)] \\ & \text{update}(\text{var } I_1, \dots, I_n; \text{st}, l) \\ & \quad := \text{st}[I_1 \mapsto (\text{var}, l, 1), \dots, I_n \mapsto (\text{var}, l, n)] \\ & \text{update}(\text{proc } I_1; K_1; \dots; I_n; K_n; \text{st}, l) \\ & \quad := \text{st}[I_1 \mapsto (\text{proc}, a_1, l, \text{size}(K_1)), \dots, I_n \mapsto (\text{proc}, a_n, l, \text{size}(K_n))] \\ & \quad \quad \text{with “fresh” addresses } a_1, \dots, a_n \\ & \quad \quad \text{where } \text{size}(D_C \ \text{var } I_1, \dots, I_n; D_P C) := n \end{aligned}$$

The Initial Symbol Table

An EPL program $P = \text{in/out } I_1, \dots, I_n; K. \in Pgm$ has a semantics of type $\mathbb{Z}^n \dashrightarrow \mathbb{Z}^n$.

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Given $(z_1, \dots, z_n) \in \mathbb{Z}^n$, we choose the **initial state**

$$s := (1, \varepsilon, \underbrace{0 : 0 : 0 : z_1 : \dots : z_n}_{\text{I/O frame}}) \in S = PC \times DS \times PS$$

The Initial Symbol Table

An EPL program $P = \text{in/out } I_1, \dots, I_n; K. \in Pgm$ has a **semantics** of type $\mathbb{Z}^n \dashrightarrow \mathbb{Z}^n$.

Given $(z_1, \dots, z_n) \in \mathbb{Z}^n$, we choose the **initial state**

$$s := (1, \varepsilon, \underbrace{0 : 0 : 0 : z_1 : \dots : z_n}_{\text{I/O frame}}) \in S = PC \times DS \times PS$$

Thus the corresponding **initial symbol table** has n entries:

$$\text{st}_{I/O}(I_j) := (\text{var}, 0, j) \quad \text{for every } j \in [n]$$

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Translation of `in/out  $I_1, \dots, I_n$ ; D C .:`

- 1 Create MAIN frame for executing C
- 2 Stop program execution after return

Translation of Programs

Translation of `in/out  $I_1, \dots, I_n; D C$` .:

- 1 Create MAIN frame for executing C
- 2 Stop program execution after return

Definition 20.2 (Translation of programs)

The mapping

$$\text{trans} : Pgm \dashrightarrow AM$$

is defined by

$$\begin{aligned} \text{trans}(\text{in/out } I_1, \dots, I_n; K.) &:= 1 : \text{CALL}(a, 0, \text{size}(K)); \\ &\quad 2 : \text{JMP}(0); \\ &\quad \text{kt}(K, \text{st}_{I/O}, a, 1) \end{aligned}$$

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Translation of D C :

- 1 Update symbol table according to D
- 2 Create code for procedures declared in D
(using the updated symbol table – recursion!)
- 3 Create code for C (using the updated symbol table)

Translation of Blocks

Translation of $D\ C$:

- 1 Update symbol table according to D
- 2 Create code for procedures declared in D
(using the updated symbol table – recursion!)
- 3 Create code for C (using the updated symbol table)

Definition 20.3 (Translation of blocks)

The mapping

$$kt : Blk \times Tab \times PC \times Lev \dashrightarrow AM$$

(“block translation”) is defined by

$$\begin{aligned} kt(D\ C, st, a, l) := & \text{dt}(D, \text{update}(D, st, l), l) \\ & \text{ct}(C, \text{update}(D, st, l), a, l) \\ & a' : \text{RET}; \end{aligned}$$

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Translation of Declarations

Translation of D :

- Generate code for the procedures declared in D

Translation of Declarations

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Definition 20.4 (Translation of declarations)

The mapping

$$dt : Dcl \times Tab \times Lev \dashrightarrow AM$$

(“declaration translation”) is defined by

$$dt(D_C \ D_V \ D_P, st, l)$$

$$:= dt(D_P, st, l)$$

$$dt(\varepsilon, st, l)$$

$$:= \varepsilon$$

$$dt(\text{proc } I_1; K_1; \dots; I_n; K_n; , st, l)$$

$$:= kt(K_1, st, a_1, l + 1)$$

$$\vdots$$

$$kt(K_n, st, a_n, l + 1)$$

where $st(I_j) = (\text{proc}, a_j, \dots, \dots)$ for every $j \in [n]$

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Translation of Commands

Definition 20.5 (Translation of commands)

The mapping

$$ct : Cmd \times Tab \times PC \times Lev \dashrightarrow AM$$

(“command translation”) is defined by

$$\begin{aligned} ct(I := A, st, a, l) &:= at(A, st, a, l) \\ &\quad a' : \text{STORE}(l - lev, off); \\ &\quad \text{if } st(I) = (\text{var}, lev, off) \\ ct(I(), st, a, l) &:= a : \text{CALL}(ca, l - lev, loc); \\ &\quad \text{if } st(I) = (\text{proc}, ca, lev, loc) \\ ct(C_1; C_2, st, a, l) &:= ct(C_1, st, a, l) \\ &\quad ct(C_2, st, a', l) \\ ct(\text{if } B \text{ then } C_1 \text{ else } C_2, st, a, l) &:= bt(B, st, a, l) \\ &\quad a' : \text{JFALSE}(a''); \\ &\quad ct(C_1, st, a' + 1, l) \\ &\quad a'' - 1 : \text{JMP}(a'''); \\ &\quad ct(C_2, st, a'', l) \\ &\quad a''' : \\ ct(\text{while } B \text{ do } C, st, a, l) &:= bt(B, st, a, l) \\ &\quad a' : \text{JFALSE}(a'' + 1); \\ &\quad ct(C, st, a' + 1, l) \\ &\quad a'' : \text{JMP}(a); \end{aligned}$$

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Translation of Boolean Expressions

Definition 20.6 (Translation of Boolean expressions)

The mapping

$$\text{bt} : BExp \times Tab \times PC \times Lev \dashrightarrow AM$$

(“Boolean expression translation”) is defined by

$$\begin{aligned} \text{bt}(A_1 < A_2, \text{st}, a, l) &:= \text{at}(A_1, \text{st}, a, l) \\ &\quad \text{at}(A_2, \text{st}, a', l) \\ &\quad a'' : \text{LT}; \end{aligned}$$

$$\begin{aligned} \text{bt}(\text{not } B, \text{st}, a, l) &:= \text{bt}(B, \text{st}, a, l) \\ &\quad a' : \text{NOT}; \end{aligned}$$

$$\begin{aligned} \text{bt}(B_1 \text{ and } B_2, \text{st}, a, l) &:= \text{bt}(B_1, \text{st}, a, l) \\ &\quad \text{bt}(B_2, \text{st}, a', l) \\ &\quad a'' : \text{AND}; \end{aligned}$$

$$\begin{aligned} \text{bt}(B_1 \text{ or } B_2, \text{st}, a, l) &:= \text{bt}(B_1, \text{st}, a, l) \\ &\quad \text{bt}(B_2, \text{st}, a', l) \\ &\quad a'' : \text{OR}; \end{aligned}$$

Translation of Arithmetic Expressions

Definition 20.7 (Translation of arithmetic expressions)

The mapping

$$at : AExp \times Tab \times PC \times Lev \dashrightarrow AM$$

(“arithmetic expression translation”) is defined by

$$at(z, st, a, l) := a : \text{LIT}(z);$$

$$at(I, st, a, l) := \begin{cases} a : \text{LIT}(z); & \text{if } st(I) = (\text{const}, z) \\ a : \text{LOAD}(l - lev, off); & \text{if } st(I) = (\text{var}, lev, off) \end{cases}$$

$$at(A_1 + A_2, st, a, l) := \begin{array}{l} at(A_1, st, a, l) \\ at(A_2, st, a', l) \\ a'' : \text{ADD}; \end{array}$$

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Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.
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Intermediate code:

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  x := y.
```

```
trans(in/out  $I_1, \dots, I_n; K.$ ) :=
```

```
1 : CALL( $a, 0, \text{size}(K)$ );  
2 : JMP(0);  
kt( $K, \text{st}_{I/O}, a, 1$ )
```

```
 $\text{st}_{I/O} = [x \mapsto (\text{var}, 0, 1)]$ 
```

Intermediate code:

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trans(in/out x;  $K.$ )
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```

```
kt( $D$   $C$ , st,  $a$ ,  $l$ ) :=  
  dt( $D$ , update( $D$ , st,  $l$ ),  $l$ )  
  ct( $C$ , update( $D$ , st,  $l$ ),  $a$ ,  $l$ )  
   $a'$  : RET;
```

$st_{I/O} = [x \mapsto (\text{var}, 0, 1)]$

Intermediate code:

```
1 : CALL( $a_0$ , 0, 1);  
2 : JMP(0);  
   kt( $K$ , stI/O,  $a_0$ , 1)
```

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```

$\text{update}(\text{var } I_1, \dots, I_n; \text{st}, l) :=$

$\text{st}[I_1 \mapsto (\text{var}, l, 1), \dots, I_n \mapsto (\text{var}, l, n)]$

$\text{update}(\text{proc } I_1; K_1; \dots; I_n; K_n; \text{st}, l) :=$

$\text{st}[I_1 \mapsto (\text{proc}, a_1, l, \text{size}(K_1)), \dots, I_n \mapsto (\text{proc}, a_n, l, \text{size}(K_n))]$

$\text{st}' = [x \mapsto (\text{var}, 0, 1),$

$y \mapsto (\text{var}, 1, 1),$

$F \mapsto (\text{proc}, a_1, 1, 0)]$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   dt(D, update(D, stI/O, 1), 1)  
   ct(C, update(D, stI/O, 1), a0, 1)  
a2 : RET;
```

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  y := 1;  
  F();  
  x := y.
```

$\text{dt}(\text{proc } I_1; K_1; \dots; I_n; K_n; \text{st}, l) :=$
 $\text{kt}(K_1, \text{st}, a_1, l + 1)$

$\vdots$
 $\text{kt}(K_n, \text{st}, a_n, l + 1)$
where $\text{st}(I_j) = (\text{proc}, a_j, \dots, \dots)$ for every $j \in [n]$

$\text{st}' = [\text{x} \mapsto (\text{var}, 0, 1),$
 $\text{y} \mapsto (\text{var}, 1, 1),$
 $\text{F} \mapsto (\text{proc}, a_1, 1, 0)]$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   dt(D, st', 1)  
   ct(C, st', a0, 1)  
a2 : RET;
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$kt(D\ C, st, a, l) :=$

```
dt(D, update(D, st, l), l)  
ct(C, update(D, st, l), a, l)  
a' : RET;
```

$st' = [x \mapsto (\text{var}, 0, 1),$
 $y \mapsto (\text{var}, 1, 1),$
 $F \mapsto (\text{proc}, a_1, 1, 0)]$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   kt(KF, st', a1, 2)  
   ct(C, st', a0, 1)  
a2 : RET;
```


Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

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in/out x;  
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```

$\text{ct}(\text{if } B \text{ then } C_1 \text{ else } C_2, \text{st}, a, l) :=$

```
bt(B, st, a, l)  
a' : JFALSE(a'');  
ct(C1, st, a' + 1, l)  
a'' - 1 : JMP(a''');  
ct(C2, st, a'', l)  
a''' :
```

$\text{st}' = [\text{x} \mapsto (\text{var}, 0, 1),$
 $\text{y} \mapsto (\text{var}, 1, 1),$
 $\text{F} \mapsto (\text{proc}, a_1, 1, 0)]$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   ct(CF, st', a1, 2)  
a3 : RET;  
   ct(C, st', a0, 1)  
a2 : RET;
```

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  x := y.
```

$\text{bt}(A_1 > A_2, \text{st}, a, l) := \text{at}(A_1, \text{st}, a, l)$
 $\text{at}(A_2, \text{st}, a', l)$
 $a'' : \text{GT};$

$\text{st}' = [\text{x} \mapsto (\text{var}, 0, 1),$
 $\text{y} \mapsto (\text{var}, 1, 1),$
 $\text{F} \mapsto (\text{proc}, a_1, 1, 0)]$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   bt(x > 1, st', a1, 2)  
a4 : JFALSE(a3);  
      ct(C1, st', a4 + 1, 2)  
a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.
```

$at(I, st, a, l) :=$

$$\begin{cases} a : LIT(z); & \text{if } st(I) = (\text{const}, z) \\ a : LOAD(l - lev, off); & \text{if } st(I) = (\text{var}, lev, off) \end{cases}$$
$$st' = [x \mapsto (\text{var}, 0, 1), \\ y \mapsto (\text{var}, 1, 1), \\ F \mapsto (\text{proc}, a_1, 1, 0)]$$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
   at(x, st', a1, 2)  
   at(1, st', a', 2)  
   GT;  
a4 : JFALSE(a3);  
      ct(C1, st', a4 + 1, 2)  
a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.  
  
 $at(z, st, a, l) := a : LIT(z);$   
  
 $st' = [x \mapsto (\text{var}, 0, 1),$   
         $y \mapsto (\text{var}, 1, 1),$   
         $F \mapsto (\text{proc}, a_1, 1, 0)]$ 
```

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
a1 : LOAD(2, 1);  
      at(1, st', a', 2)  
      GT;  
a4 : JFALSE(a3);  
      ct(C1, st', a4 + 1, 2)  
a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.
```

```
ct( $I := A, st, a, l$ ) :=  
  at( $A, st, a, l$ )  
   $a' : \text{STORE}(l - lev, off)$ ;  
if  $st(I) = (\text{var}, lev, off)$   
 $st' = [x \mapsto (\text{var}, 0, 1),$   
   $y \mapsto (\text{var}, 1, 1),$   
   $F \mapsto (\text{proc}, a_1, 1, 0)]$ 
```

Intermediate code:

```
1 : CALL( $a_0, 0, 1$ );  
2 : JMP(0);  
 $a_1$  : LOAD(2, 1);  
      LIT(1);  
      GT;  
 $a_4$  : JFALSE( $a_3$ );  
      ct( $C_1, st', a_4 + 1, 2$ )  
 $a_3$  : RET;  
      ct( $C, st', a_0, 1$ )  
 $a_2$  : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.  
  
at( $A_1 + A_2$ , st, a, l) := at( $A_1$ , st, a, l)  
                        at( $A_2$ , st, a', l)  
                        a'' : ADD;  
  
st' = [x  $\mapsto$  (var, 0, 1),  
       y  $\mapsto$  (var, 1, 1),  
       F  $\mapsto$  (proc, a1, 0)]
```

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
a1 : LOAD(2, 1);  
      LIT(1);  
      GT;  
a4 : JFALSE(a3);  
      at(y * x, st', a4 + 1, 2)  
      STORE(1, 1);  
      at(x - 1, st', a', 2)  
      STORE(2, 1);  
      ct(F(), st', a'', 2)  
      a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.  
  
ct(I(), st, a, l) :=  
  a : CALL(ca, l - lev, loc);  
if st(I) = (proc, ca, lev, loc)  
st' = [x ↦ (var, 0, 1),  
       y ↦ (var, 1, 1),  
       F ↦ (proc, a1, 1, 0)]
```

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
a1 : LOAD(2, 1);  
      LIT(1);  
      GT;  
a4 : JFALSE(a3);  
      LOAD(1, 1);  
      LOAD(2, 1);  
      MULT;  
      STORE(1, 1);  
      LOAD(2, 1);  
      LIT(1);  
      SUB;  
      STORE(2, 1);  
      ct(F(), st', a'', 2)  
a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```

Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.
```

$$st' = [x \mapsto (\text{var}, 0, 1), \\ y \mapsto (\text{var}, 1, 1), \\ F \mapsto (\text{proc}, a_1, 1, 0)]$$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
a1 : LOAD(2, 1);  
      LIT(1);  
      GT;  
a4 : JFALSE(a3);  
      LOAD(1, 1);  
      LOAD(2, 1);  
      MULT;  
      STORE(1, 1);  
      LOAD(2, 1);  
      LIT(1);  
      SUB;  
      STORE(2, 1);  
      CALL(a1, 1, 0);  
a3 : RET;  
      ct(C, st', a0, 1)  
a2 : RET;
```


Example: Factorial Function I

Example 20.8 (Factorial function (cf. Ex. 19.1))

Source code:

```
in/out x;  
var y;  
proc F;  
  if x > 1 then  
    y := y * x;  
    x := x - 1;  
    F()  
  y := 1;  
  F();  
  x := y.
```

$$st' = [x \mapsto (\text{var}, 0, 1), \\ y \mapsto (\text{var}, 1, 1), \\ F \mapsto (\text{proc}, a_1, 1, 0)]$$

Intermediate code:

```
1 : CALL(a0, 0, 1);  
2 : JMP(0);  
a1 : LOAD(2, 1);  
      LIT(1);  
      GT;  
a4 : JFALSE(a3);  
      LOAD(1, 1);  
      LOAD(2, 1);  
      MULT;  
      STORE(1, 1);  
      LOAD(2, 1);  
      LIT(1);  
      SUB;  
      STORE(2, 1);  
      CALL(a1, 1, 0);  
a3 : RET;  
a0 : LIT(1);  
      STORE(0, 1);  
      CALL(a1, 0, 0);  
      LOAD(0, 1);  
      STORE(1, 1);  
a2 : RET;
```

Example: Factorial Function II

Example 20.8 (Factorial function; continued)

Code with symbolic addresses:

```
1 : CALL(a0,0,1);
2 : JMP(0);
a1 : LOAD(2,1);
    LIT(1);
    GT;
a4 : JFALSE(a3);
    LOAD(1,1);
    LOAD(2,1);
    MULT;
    STORE(1,1);
    LOAD(2,1);
    LIT(1);
    SUB;
    STORE(2,1);
    CALL(a1,1,0);
a3 : RET;
a0 : LIT(1);
    STORE(0,1);
    CALL(a1,0,0);
    LOAD(0,1);
    STORE(1,1);
a2 : RET;
```

Linearized

($a_0 = 17, a_1 = 3, a_2 = 22, a_3 = 16, a_4 = 6$):

```
1 : CALL(17,0,1);
2 : JMP(0);
3 : LOAD(2,1);
4 : LIT(1);
5 : GT;
6 : JFALSE(16);
7 : LOAD(1,1);
8 : LOAD(2,1);
9 : MULT;
10 : STORE(1,1);
11 : LOAD(2,1);
12 : LIT(1);
13 : SUB;
14 : STORE(2,1);
15 : CALL(3,1,0);
16 : RET;
17 : LIT(1);
18 : STORE(0,1);
19 : CALL(3,0,0);
20 : LOAD(0,1);
21 : STORE(1,1);
22 : RET;
```

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

<i>PC</i>	<i>DS</i>	<i>PS</i>
1	ε	<u>0 : 0 : 0 : 2</u>

```
1 : CALL(17,0,1);
2 : JMP(0);
3 : LOAD(2,1);
4 : LIT(1);
5 : GT;
6 : JFALSE(16);
7 : LOAD(1,1);
8 : LOAD(2,1);
9 : MULT;
10 : STORE(1,1);
11 : LOAD(2,1);
12 : LIT(1);
13 : SUB;
14 : STORE(2,1);
15 : CALL(3,1,0);
16 : RET;
17 : LIT(1);
18 : STORE(0,1);
19 : CALL(3,0,0);
20 : LOAD(0,1);
21 : STORE(1,1);
22 : RET;
```

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);			
4: LIT(1);			
5: GT;			
6: JFALSE(16);			
7: LOAD(1,1);			
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);			
5: GT;			
6: JFALSE(16);			
7: LOAD(1,1);			
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;			
6: JFALSE(16);			
7: LOAD(1,1);			
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);			
7: LOAD(1,1);			
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20:4:3:2:1:0:0:0:2</u>
7: LOAD(1,1);			
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20:4:3:2:1:0:0:0:2</u>
7: LOAD(1,1);	5	2:1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
8: LOAD(2,1);			
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20:4:3:2:1:0:0:0:2</u>
7: LOAD(1,1);	5	2:1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
8: LOAD(2,1);	6	1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
9: MULT;			
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20:4:3:2:1:0:0:0:2</u>
7: LOAD(1,1);	5	2:1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
8: LOAD(2,1);	6	1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
9: MULT;	7	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
10: STORE(1,1);			
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	<u>4:3:2:0:0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0:0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1:0:0:0:2</u>
5: GT;	3	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20:4:3:2:1:0:0:0:2</u>
7: LOAD(1,1);	5	2:1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
8: LOAD(2,1);	6	1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
9: MULT;	7	ε	<u>3:2:20:4:3:2:1:0:0:0:2</u>
10: STORE(1,1);	8	1	<u>3:2:20:4:3:2:1:0:0:0:2</u>
11: LOAD(2,1);			
12: LIT(1);			
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

<i>PC</i>	<i>DS</i>	<i>PS</i>
1	ε	<u>0 : 0 : 0 : 2</u>
17	ε	<u>4 : 3 : 2 : 0 : 0 : 0 : 0 : 2</u>
18	1	<u>4 : 3 : 2 : 0 : 0 : 0 : 0 : 2</u>
19	ε	<u>4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
3	ε	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
4	2	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
5	2 : 1	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
6	1	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
7	ε	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
8	1	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>
9	1 : 2	<u>3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2</u>

```

1 : CALL(17,0,1);
2 : JMP(0);
3 : LOAD(2,1);
4 : LIT(1);
5 : GT;
6 : JFALSE(16);
7 : LOAD(1,1);
8 : LOAD(2,1);
9 : MULT;
10 : STORE(1,1);
11 : LOAD(2,1);
12 : LIT(1);
13 : SUB;
14 : STORE(2,1);
15 : CALL(3,1,0);
16 : RET;
17 : LIT(1);
18 : STORE(0,1);
19 : CALL(3,0,0);
20 : LOAD(0,1);
21 : STORE(1,1);
22 : RET;
    
```

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	$0:0:0:2$
2: JMP(0);	17	ε	$4:3:2:0:0:0:0:2$
3: LOAD(2,1);	18	1	$4:3:2:0:0:0:0:2$
4: LIT(1);	19	ε	$4:3:2:1:0:0:0:2$
5: GT;	3	ε	$3:2:20:4:3:2:1:0:0:0:2$
6: JFALSE(16);	4	2	$3:2:20:4:3:2:1:0:0:0:2$
7: LOAD(1,1);	5	2:1	$3:2:20:4:3:2:1:0:0:0:2$
8: LOAD(2,1);	6	1	$3:2:20:4:3:2:1:0:0:0:2$
9: MULT;	7	ε	$3:2:20:4:3:2:1:0:0:0:2$
10: STORE(1,1);	8	1	$3:2:20:4:3:2:1:0:0:0:2$
11: LOAD(2,1);	9	1:2	$3:2:20:4:3:2:1:0:0:0:2$
12: LIT(1);	10	2	$3:2:20:4:3:2:1:0:0:0:2$
13: SUB;			
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	$0:0:0:2$
2: JMP(0);	17	ε	$4:3:2:0:0:0:0:2$
3: LOAD(2,1);	18	1	$4:3:2:0:0:0:0:2$
4: LIT(1);	19	ε	$4:3:2:1:0:0:0:2$
5: GT;	3	ε	$3:2:20:4:3:2:1:0:0:0:2$
6: JFALSE(16);	4	2	$3:2:20:4:3:2:1:0:0:0:2$
7: LOAD(1,1);	5	2:1	$3:2:20:4:3:2:1:0:0:0:2$
8: LOAD(2,1);	6	1	$3:2:20:4:3:2:1:0:0:0:2$
9: MULT;	7	ε	$3:2:20:4:3:2:1:0:0:0:2$
10: STORE(1,1);	8	1	$3:2:20:4:3:2:1:0:0:0:2$
11: LOAD(2,1);	9	1:2	$3:2:20:4:3:2:1:0:0:0:2$
12: LIT(1);	10	2	$3:2:20:4:3:2:1:0:0:0:2$
13: SUB;	11	ε	$3:2:20:4:3:2:2:0:0:0:2$
14: STORE(2,1);			
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	<u>0:0:0:2</u>
2: JMP(0);	17	ε	4:3:2:0: <u>0:0:0:2</u>
3: LOAD(2,1);	18	1	<u>4:3:2:0</u> : <u>0:0:0:2</u>
4: LIT(1);	19	ε	<u>4:3:2:1</u> : <u>0:0:0:2</u>
5: GT;	3	ε	3:2:20: <u>4:3:2:1</u> : <u>0:0:0:2</u>
6: JFALSE(16);	4	2	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
7: LOAD(1,1);	5	2:1	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
8: LOAD(2,1);	6	1	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
9: MULT;	7	ε	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
10: STORE(1,1);	8	1	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
11: LOAD(2,1);	9	1:2	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
12: LIT(1);	10	2	<u>3:2:20</u> : <u>4:3:2:1</u> : <u>0:0:0:2</u>
13: SUB;	11	ε	<u>3:2:20</u> : <u>4:3:2:2</u> : <u>0:0:0:2</u>
14: STORE(2,1);	12	2	<u>3:2:20</u> : <u>4:3:2:2</u> : <u>0:0:0:2</u>
15: CALL(3,1,0);			
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;			
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;	14	1	3:2:20:4:3:2:2:0:0:0:2
17: LIT(1);			
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;	14	1	3:2:20:4:3:2:2:0:0:0:2
17: LIT(1);	15	ε	3:2:20:4:3:2:2:0:0:0:1
18: STORE(0,1);			
19: CALL(3,0,0);			
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1 : CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2 : JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3 : LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4 : LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5 : GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6 : JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7 : LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8 : LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9 : MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10 : STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11 : LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12 : LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13 : SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14 : STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15 : CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16 : RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17 : LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18 : STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19 : CALL(3,0,0);			
20 : LOAD(0,1);			
21 : STORE(1,1);			
22 : RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;	14	1	3:2:20:4:3:2:2:0:0:0:2
17: LIT(1);	15	ε	3:2:20:4:3:2:2:0:0:0:1
18: STORE(0,1);	3	ε	6:2:16:3:2:20:4:3:2:2:0:0:0:1
19: CALL(3,0,0);	4	1	6:2:16:3:2:20:4:3:2:2:0:0:0:1
20: LOAD(0,1);			
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	$0:0:0:2$
2: JMP(0);	17	ε	$4:3:2:0:0:0:0:2$
3: LOAD(2,1);	18	1	$4:3:2:0:0:0:0:2$
4: LIT(1);	19	ε	$4:3:2:1:0:0:0:2$
5: GT;	3	ε	$3:2:20:4:3:2:1:0:0:0:2$
6: JFALSE(16);	4	2	$3:2:20:4:3:2:1:0:0:0:2$
7: LOAD(1,1);	5	2:1	$3:2:20:4:3:2:1:0:0:0:2$
8: LOAD(2,1);	6	1	$3:2:20:4:3:2:1:0:0:0:2$
9: MULT;	7	ε	$3:2:20:4:3:2:1:0:0:0:2$
10: STORE(1,1);	8	1	$3:2:20:4:3:2:1:0:0:0:2$
11: LOAD(2,1);	9	1:2	$3:2:20:4:3:2:1:0:0:0:2$
12: LIT(1);	10	2	$3:2:20:4:3:2:1:0:0:0:2$
13: SUB;	11	ε	$3:2:20:4:3:2:2:0:0:0:2$
14: STORE(2,1);	12	2	$3:2:20:4:3:2:2:0:0:0:2$
15: CALL(3,1,0);	13	2:1	$3:2:20:4:3:2:2:0:0:0:2$
16: RET;	14	1	$3:2:20:4:3:2:2:0:0:0:2$
17: LIT(1);	15	ε	$3:2:20:4:3:2:2:0:0:0:1$
18: STORE(0,1);	3	ε	$6:2:16:3:2:20:4:3:2:2:0:0:0:1$
19: CALL(3,0,0);	4	1	$6:2:16:3:2:20:4:3:2:2:0:0:0:1$
20: LOAD(0,1);	5	1:1	$6:2:16:3:2:20:4:3:2:2:0:0:0:1$
21: STORE(1,1);			
22: RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:

	<i>PC</i>	<i>DS</i>	<i>PS</i>
1 : CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2 : JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3 : LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4 : LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5 : GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6 : JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7 : LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8 : LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9 : MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10 : STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11 : LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12 : LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13 : SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14 : STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15 : CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16 : RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17 : LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18 : STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19 : CALL(3,0,0);	4	1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
20 : LOAD(0,1);	5	1 : 1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
21 : STORE(1,1);	6	0	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
22 : RET;			

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2: JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3: LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4: LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5: GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6: JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7: LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8: LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9: MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10: STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11: LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12: LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13: SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14: STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15: CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16: RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17: LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18: STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19: CALL(3,0,0);	4	1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
20: LOAD(0,1);	5	1 : 1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
21: STORE(1,1);	6	0	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
22: RET;	16	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;	14	1	3:2:20:4:3:2:2:0:0:0:2
17: LIT(1);	15	ε	3:2:20:4:3:2:2:0:0:0:1
18: STORE(0,1);	3	ε	6:2:16:3:2:20:4:3:2:2:0:0:0:1
19: CALL(3,0,0);	4	1	6:2:16:3:2:20:4:3:2:2:0:0:0:1
20: LOAD(0,1);	5	1:1	6:2:16:3:2:20:4:3:2:2:0:0:0:1
21: STORE(1,1);	6	0	6:2:16:3:2:20:4:3:2:2:0:0:0:1
22: RET;	16	ε	6:2:16:3:2:20:4:3:2:2:0:0:0:1

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2: JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3: LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4: LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5: GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6: JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7: LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8: LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9: MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10: STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11: LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12: LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13: SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14: STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15: CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16: RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17: LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18: STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19: CALL(3,0,0);	4	1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
20: LOAD(0,1);	5	1 : 1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
21: STORE(1,1);	6	0	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
22: RET;	16	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	16	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	20	ε	4 : 3 : 2 : 2 : 0 : 0 : 0 : 1

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1 : CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2 : JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3 : LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4 : LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5 : GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6 : JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7 : LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8 : LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9 : MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10 : STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11 : LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12 : LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13 : SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14 : STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15 : CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16 : RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17 : LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18 : STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19 : CALL(3,0,0);	4	1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
20 : LOAD(0,1);	5	1 : 1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
21 : STORE(1,1);	6	0	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
22 : RET;	16	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	16	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	20	ε	4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	21	2	4 : 3 : 2 : 2 : 0 : 0 : 0 : 1

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0 : 0 : 0 : 2
2: JMP(0);	17	ε	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
3: LOAD(2,1);	18	1	4 : 3 : 2 : 0 : 0 : 0 : 0 : 2
4: LIT(1);	19	ε	4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
5: GT;	3	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
6: JFALSE(16);	4	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
7: LOAD(1,1);	5	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
8: LOAD(2,1);	6	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
9: MULT;	7	ε	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
10: STORE(1,1);	8	1	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
11: LOAD(2,1);	9	1 : 2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
12: LIT(1);	10	2	3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2
13: SUB;	11	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
14: STORE(2,1);	12	2	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
15: CALL(3,1,0);	13	2 : 1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
16: RET;	14	1	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2
17: LIT(1);	15	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
18: STORE(0,1);	3	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
19: CALL(3,0,0);	4	1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
20: LOAD(0,1);	5	1 : 1	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
21: STORE(1,1);	6	0	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
22: RET;	16	ε	6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	16	ε	3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	20	ε	4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	21	2	4 : 3 : 2 : 2 : 0 : 0 : 0 : 1
	22	ε	4 : 3 : 2 : 2 : 0 : 0 : 0 : 2

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1 : CALL(17,0,1);	1	ε	$0 : 0 : 0 : 2$
2 : JMP(0);	17	ε	$4 : 3 : 2 : 0 : 0 : 0 : 0 : 2$
3 : LOAD(2,1);	18	1	$4 : 3 : 2 : 0 : 0 : 0 : 0 : 2$
4 : LIT(1);	19	ε	$4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
5 : GT;	3	ε	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
6 : JFALSE(16);	4	2	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
7 : LOAD(1,1);	5	$2 : 1$	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
8 : LOAD(2,1);	6	1	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
9 : MULT;	7	ε	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
10 : STORE(1,1);	8	1	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
11 : LOAD(2,1);	9	$1 : 2$	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
12 : LIT(1);	10	2	$3 : 2 : 20 : 4 : 3 : 2 : 1 : 0 : 0 : 0 : 2$
13 : SUB;	11	ε	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2$
14 : STORE(2,1);	12	2	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2$
15 : CALL(3,1,0);	13	$2 : 1$	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2$
16 : RET;	14	1	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 2$
17 : LIT(1);	15	ε	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
18 : STORE(0,1);	3	ε	$6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
19 : CALL(3,0,0);	4	1	$6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
20 : LOAD(0,1);	5	$1 : 1$	$6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
21 : STORE(1,1);	6	0	$6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
22 : RET;	16	ε	$6 : 2 : 16 : 3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
	16	ε	$3 : 2 : 20 : 4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
	20	ε	$4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
	21	2	$4 : 3 : 2 : 2 : 0 : 0 : 0 : 1$
	22	ε	$4 : 3 : 2 : 2 : 0 : 0 : 0 : 2$
	2	ε	$0 : 0 : 0 : 2$

Example: Factorial Function III

Example 20.8 (Factorial function; continued)

Computation for $x = 2$:	<i>PC</i>	<i>DS</i>	<i>PS</i>
1: CALL(17,0,1);	1	ε	0:0:0:2
2: JMP(0);	17	ε	4:3:2:0:0:0:0:2
3: LOAD(2,1);	18	1	4:3:2:0:0:0:0:2
4: LIT(1);	19	ε	4:3:2:1:0:0:0:2
5: GT;	3	ε	3:2:20:4:3:2:1:0:0:0:2
6: JFALSE(16);	4	2	3:2:20:4:3:2:1:0:0:0:2
7: LOAD(1,1);	5	2:1	3:2:20:4:3:2:1:0:0:0:2
8: LOAD(2,1);	6	1	3:2:20:4:3:2:1:0:0:0:2
9: MULT;	7	ε	3:2:20:4:3:2:1:0:0:0:2
10: STORE(1,1);	8	1	3:2:20:4:3:2:1:0:0:0:2
11: LOAD(2,1);	9	1:2	3:2:20:4:3:2:1:0:0:0:2
12: LIT(1);	10	2	3:2:20:4:3:2:1:0:0:0:2
13: SUB;	11	ε	3:2:20:4:3:2:2:0:0:0:2
14: STORE(2,1);	12	2	3:2:20:4:3:2:2:0:0:0:2
15: CALL(3,1,0);	13	2:1	3:2:20:4:3:2:2:0:0:0:2
16: RET;	14	1	3:2:20:4:3:2:2:0:0:0:2
17: LIT(1);	15	ε	3:2:20:4:3:2:2:0:0:0:1
18: STORE(0,1);	3	ε	6:2:16:3:2:20:4:3:2:2:0:0:0:1
19: CALL(3,0,0);	4	1	6:2:16:3:2:20:4:3:2:2:0:0:0:1
20: LOAD(0,1);	5	1:1	6:2:16:3:2:20:4:3:2:2:0:0:0:1
21: STORE(1,1);	6	0	6:2:16:3:2:20:4:3:2:2:0:0:0:1
22: RET;	16	ε	6:2:16:3:2:20:4:3:2:2:0:0:0:1
	16	ε	3:2:20:4:3:2:2:0:0:0:1
	20	ε	4:3:2:2:0:0:0:1
	21	2	4:3:2:2:0:0:0:1
	22	ε	4:3:2:2:0:0:0:2
	2	ε	0:0:0:2
	0	ε	0:0:0:2

- 1 Repetition: Intermediate Code
- 2 The Symbol Table
- 3 Translation of Programs
- 4 Translation of Blocks
- 5 Translation of Declarations
- 6 Translation of Commands
- 7 Translation of Expressions
- 8 A Translation Example
- 9 Correctness of the Translation
- 10 Outlook

Theorem 20.9 (Correctness of translation)

For every $P \in Pgm$, $n \in \mathbb{N}$, and $(z_1, \dots, z_n), (z'_1, \dots, z'_n) \in \mathbb{Z}^n$:

$$\begin{aligned} & \mathfrak{M}[[P]](z_1, \dots, z_n) = (z'_1, \dots, z'_n) \\ \iff & \llbracket \text{trans}(P) \rrbracket(1, \varepsilon, 0 : 0 : 0 : z_1 : \dots : z_n) = (0, \varepsilon, 0 : 0 : 0 : z'_1 : \dots : z'_n) \end{aligned}$$

Theorem 20.9 (Correctness of translation)

For every $P \in Pgm$, $n \in \mathbb{N}$, and $(z_1, \dots, z_n), (z'_1, \dots, z'_n) \in \mathbb{Z}^n$:

$$\begin{aligned} & \mathfrak{M}[[P]](z_1, \dots, z_n) = (z'_1, \dots, z'_n) \\ \iff & [[\text{trans}(P)]](1, \varepsilon, 0 : 0 : 0 : z_1 : \dots : z_n) = (0, \varepsilon, 0 : 0 : 0 : z'_1 : \dots : z'_n) \end{aligned}$$

Proof.

see M. Mohnen: *A Compiler Correctness Proof for the Static Link Technique by means of Evolving Algebras*, Fundamenta Informaticae 29(3), 1997, pp. 257–303 □

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- 10 Outlook

- ❶ Error handling in (top-down) parsing
- ❷ Strong noncircularity of attribute grammars
- ❸ More about code generation
 - Static & dynamic data structures
 - Compiler backend (register allocation, ...)
- ❹ Code analysis & optimization