

# Concurrency Theory WS 2013/2014

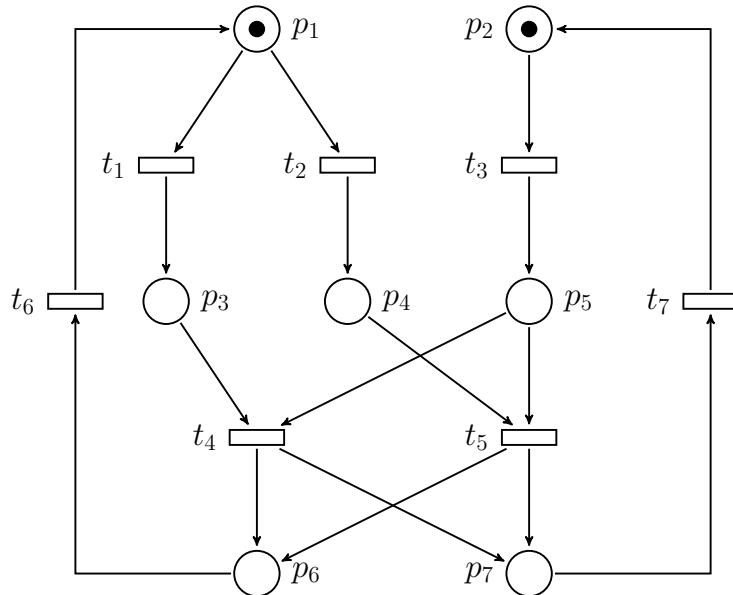
## — Series 10 —

Hand in until January 21st before the exercise class.

---

### Exercise 1 (Marking Graph and Distributed Runs)

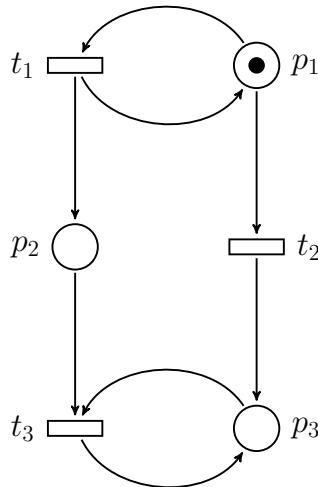
(3 Points)



- 1) Give the marking graph of the Petri net shown above.
- 2) Give at least three distributed runs which cover at least four transition of this Petri net. If the Petri net admits an infinite distributed run, please provide it.

### Exercise 2 (Infinite Marking Graph)

(3 Points)



- 1) Give the (partial) marking graph of the Petri net shown above and show that it is infinite.
- 2) Since the marking graphs of a Petri net can be infinite, *coverability tree* (or coverability graph) is defined for Petri net, which is always finite. In order to represent *infinite* tokens in a place, a special symbol  $\omega$  is introduced, and for all natural number  $n$  it has the following properties:  $\omega > n$ ,  $\omega \pm n = \omega$  and  $\omega \geq \omega$ .

The coverability tree for a Petri net  $(P, T, F, M_0)$  can then be computed by following algorithm.

```

s.1) Label  $M_0$  as the root and tag it "new".
s.2) While "new" markings exist, do the following:
    s.2.1) Select a new marking  $M$ .
    s.2.2) If  $M$  is identical to a marking on the path from the root
        to  $M$ , tag  $M$  "old" and to another new marking.
    s.2.3) If no transitions are enabled at  $M$ , tag  $M$  as "dead-end".
    s.2.4) While there exist enabled transitions at  $M$ , do the
        following for each enabled transition  $t$  at  $M$ :
        s.2.4.1) Obtain the marking  $M'$  that results from firing  $t$  at  $M$ .
        s.2.4.2) On the path from the root to  $M$ , if there exist a marking
             $M''$  such that  $M'(p) \geq M''(p)$  for each place  $p$  and
             $M' \neq M''$ , replace  $M'(p)$  by  $\omega$  for each  $p$ 
            such that  $M'(p) > M''(p)$ .
        s.2.4.4) Introduce  $M'$  as a node, draw an arc with label  $t$ 
            from  $M$  to  $M'$ , and tag  $M'$  "new".
    
```

Compute the coverability tree of the above Petri net.

- 3) The *synchronic distance* between two transitions  $t$  and  $t'$  in an elementary system net  $(P, T, F, M_0)$  is defined by

$$d_{t,t'} = \max_{\sigma} |\#_{\sigma}(t) - \#_{\sigma}(t')|$$

where  $\sigma$  is a sequential run starting at any reachable marking  $M$  of  $M_0$  and  $\#_{\sigma}(t)$  is the number of times  $t$  fires in  $\sigma$ . Please give  $d_{t_1,t_2}$ ,  $d_{t_1,t_3}$  and  $d_{t_2,t_3}$  for the above net.

### Exercise 3 (Modelling with Petri nets)

(4 Points)

- 1) Model the traffic light system at a cross road in *Germany* by means of elementary Petri nets. If you don't know how it works, you should do some on-the-spot investigation.
- 2) In Exercise Series 2, we have introduced Hyman's algorithm for mutual exclusion. Please model this algorithm by means of elementary Petri nets.