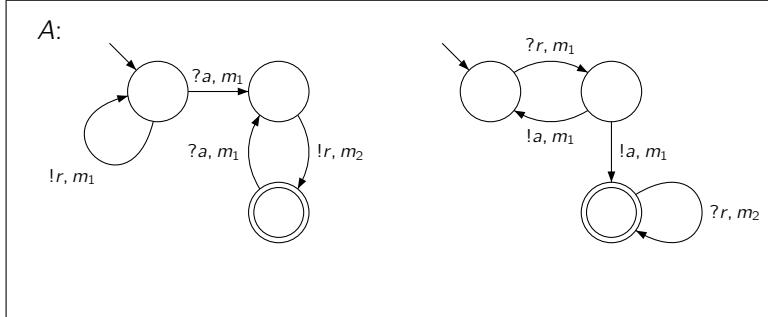


**Exercise 1 (Safe Realizability):**

**(6 Points)**

Given the following CFM A:



Show that A is not safe by finding a configuration that is reachable from the initial configuration  $\gamma_0$  of A and from which a final configuration  $\gamma_e$  cannot be obtained. Justify your answer by indicating the sequence of configurations leading from the initial configuration  $\gamma_0$  to the deadlock configuration  $\gamma_d$  and arguing why a final configuration is not reachable from  $\gamma_d$ .

**Exercise 2 (Deadlock freeness):**

**(7 Points)**

For well-formed language  $L \subseteq Act^*$ , and proper word  $w \in Act^*$ , i.e.,  $w$  is a prefix of a well-formed word, let:  $L \models^{df} w$  iff  $(\forall p \in \mathcal{P}. \exists v \in L. w \upharpoonright p \text{ is a prefix of } v \upharpoonright p)$ . Language  $L$  is closed under  $\models^{df}$  iff  $L \models^{df} w$  implies  $w \in Pref(L)$ .

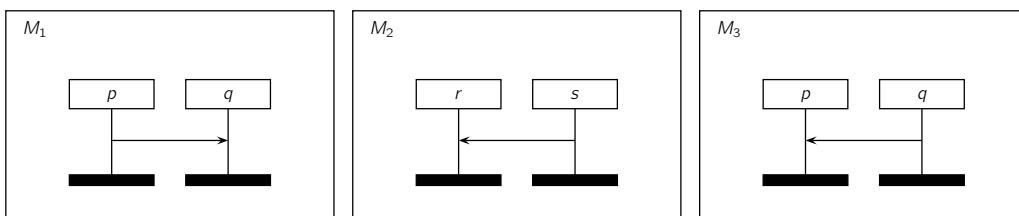
A language  $L \subseteq Act^*$  is closed under  $\models^{df}$  iff for all  $v, w \in Pref(L)$  and all processes  $p \in \mathcal{P}$ :  $(v \upharpoonright p = w \upharpoonright p \text{ and } vx \in Pref(L) \text{ for } x \in Act_p \text{ and } wx \text{ is prefix of a well-formed word})$  implies  $wx \in Pref(L)$ .

Prove the following statement: A language  $L$  is closed under  $\models^{df}$  iff  $L$  is closed under  $\models^{df}$ .

**Exercise 3 (Safe Realizability):**

**(6 Points)**

Check (i.e., by using the definitions) whether language  $L_i$  ( $i \in \{1, 2\}$ ) is closed under  $\models$  and  $\models^{df}$ :



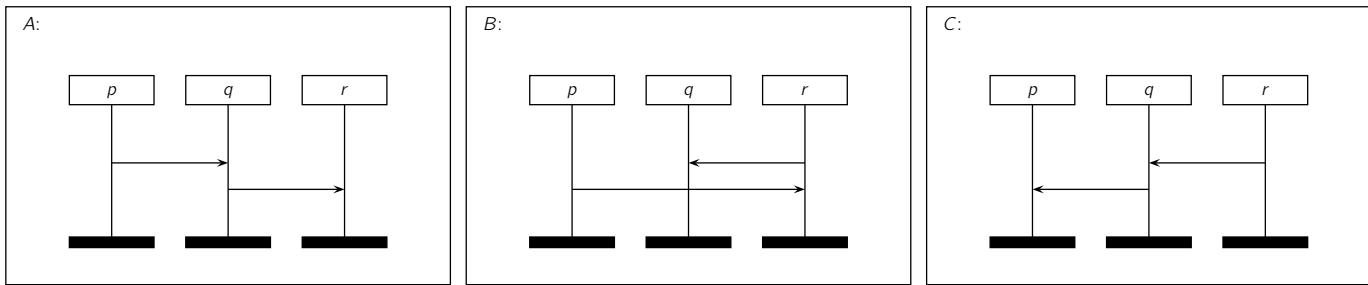
$$L_1 = \{w \mid w \in Lin(\{M_1, M_2\})\} \text{ and } L_2 = \{w \mid w \in Lin(\{M_1, M_3\})\}$$

Which of the languages is realizable or even safely realizable? Justify your answers.

(Note that  $Lin(\{M, M'\}) = Lin(M \cdot M') \cup Lin(M' \cdot M)$ .)

**Exercise 4 (Regular Expressions and Realizability):**

**(6 Points)**



Check whether the following regular expressions are realisable or not:

- $\alpha_1 = A^* + B^* + C^*$ ,
- $\alpha_2 = A \cdot (B + C)^*$ ,
- $\alpha_3 = (A \cdot B \cdot C)^*$ ,
- $\alpha_4 = (A \cdot B)^*$ ,
- $\alpha_5 = (A \cdot C)^*$ ,

and if so, whether they can be realized by a universally bounded, deadlock-free CFM?