

## Introduction to Model Checking Winter term 2011/2012

### – Series 10 –

Hand in on January 11<sup>th</sup> before the exercise class.

#### Exercise 1

(2 points)

Provide a non-deterministic Büchi Automaton  $\mathcal{A}_i$  for each of the following LTL-formulas  $\varphi_i$  such that  $\mathcal{L}_\omega(\mathcal{A}_i) = \text{Words}(\varphi_i)$ :

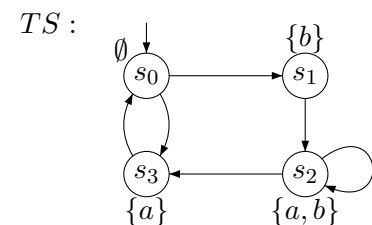
- (a)  $\varphi_1 = \Box(a \vee \neg \bigcirc b)$
- (b)  $\varphi_2 = \Diamond a \vee \Box \Diamond(a \leftrightarrow b)$
- (c)  $\varphi_3 = \bigcirc \bigcirc (a \vee \Diamond \Box b)$

*Hint:* You can infer the NBA directly, i.e. without the algorithm given in the lecture.

#### Exercise 2

(6 points)

We consider the LTL formula  $\varphi = \Box(a \rightarrow (\neg b \mathbf{U}(a \wedge b)))$  over the set  $AP = \{a, b\}$  of atomic propositions and want to check  $TS \models \varphi$  wrt. the transition system outlined on the right.



- a) To check  $TS \models \varphi$ , convert  $\neg \varphi$  into an equivalent LTL-formula  $\psi$  which is constructed according to the following grammar:

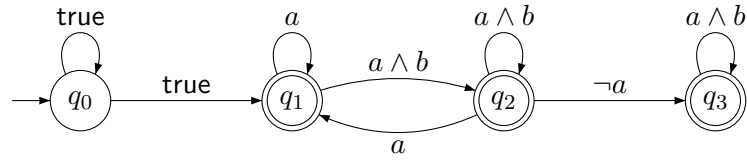
$$\varphi ::= \text{true} \mid \text{false} \mid a \mid b \mid \varphi \wedge \varphi \mid \neg \varphi \mid \bigcirc \varphi \mid \varphi \mathbf{U} \varphi.$$

Then construct  $\text{closure}(\psi)$ .

- b) Give the elementary sets wrt.  $\text{closure}(\psi)$ !
- c) Construct the GNBA  $\mathcal{G}_\psi$  by providing its initial states, its acceptance set and its transition relation. Use the algorithm given in the lecture.  
*Hint: It suffices to provide the transition relation as a table.*
- d) Now, construct an NBA  $\mathcal{A}_{\neg \varphi}$  **directly** from  $\neg \varphi$ , i.e. without relying on  $\mathcal{G}_\psi$ .  
*Hint: Four states suffice!*
- e) Construct  $TS \otimes \mathcal{A}_{\neg \varphi}$ .
- f) Use the Nested DFS algorithm from the lecture to check  $TS \models \varphi$ . Therefore sketch the algorithm's main steps and interpret its outcome!

**Exercise 3****(2 points)**

Consider the GNBA  $\mathcal{G}$  over the alphabet  $\Sigma = 2^{\{a,b\}}$  and the set  $\mathcal{F} = \{\{q_1, q_3\}, \{q_2\}\}$  of acceptance sets:



- a) Provide an LTL formula  $\varphi$  such that  $Word(\varphi) = \mathcal{L}_\omega(\mathcal{G})$ . Justify your answer!
- b) Depict an NBA  $\mathcal{A}$  with  $\mathcal{L}_\omega(\mathcal{A}) = \mathcal{L}_\omega(\mathcal{G})$ .