

## Introduction to Model Checking

### Winter term 2013/2014

#### – Series 7 –

Hand in on December 11<sup>th</sup> before the exercise class or in the box in front of the chair's secretary's office.

#### Exercise 1

(2 points)

- a) Give an LT property that is not an  $\omega$ -regular property.
- b) Give an  $\omega$ -regular property that is not expressible in LTL.

#### Exercise 2

(3 points)

Let  $\varphi$  and  $\psi$  be LTL formulae. Consider the following new operators:

- a) “At next”  $\varphi \mathbf{AX} \psi$ :

$$\begin{aligned}
 A_1 A_2 \dots \models \varphi \mathbf{AX} \psi &\iff \text{for all } i \geq 0 \text{ where } A_i A_{i+1} \dots \models \psi, \\
 &\quad \text{for which there exists no } 0 \leq j < i \text{ where } A_j A_{j+1} \dots \models \psi, \\
 &\quad A_i A_{i+1} \dots \models \varphi \text{ holds}
 \end{aligned}$$

- b) “While”  $\varphi \mathbf{WH} \psi$ :

$$\begin{aligned}
 A_1 A_2 \dots \models \varphi \mathbf{WH} \psi &\iff \text{for all } i \geq 0 \text{ where } A_j A_{j+1} \dots \models \psi \text{ for all } 0 \leq j < i, \\
 &\quad A_k A_{k+1} \dots \models \varphi \text{ for all } 0 \leq k < i
 \end{aligned}$$

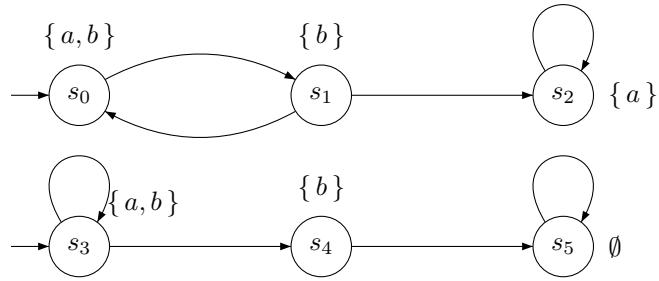
- c) “Before”  $\varphi \mathbf{B} \psi$ :

$$\begin{aligned}
 A_1 A_2 \dots \models \varphi \mathbf{B} \psi &\iff \text{for all } i \geq 0 \text{ where } A_i A_{i+1} \dots \models \psi, \\
 &\quad \text{there exists some } 0 \leq j < i \text{ where } A_j A_{j+1} \dots \models \varphi
 \end{aligned}$$

Show that these operators are LTL-definable by providing equivalent LTL formulae. You may use both the until and weak until operator.

### Exercise 3

(4 points)



Consider the transition system  $T$  above with the set  $AP = \{a, b, c\}$  of atomic propositions. Note that this is a single transition system with two initial states. Consider the LTL fairness assumption

$$fair = (\square \diamond (a \ \& \ b) \rightarrow \square \diamond \neg c) \ \wedge \ (\diamond \square (a \ \& \ b) \rightarrow \square \diamond \neg b).$$

Questions:

- Determine the fair paths in  $T$ , i.e., the initial, infinite paths satisfying  $fair$
- For each of the following LTL formulae:

$$\begin{aligned}\varphi_1 &= b \mathsf{U} \ \square \neg b \\ \varphi_2 &= b \mathsf{W} \ \square \neg b \\ \varphi_3 &= \bigcirc \bigcirc b \mathsf{U} \ \square \neg b\end{aligned}$$

determine whether  $T \models_{fair} \varphi_i$ . In case  $T \not\models_{fair} \varphi_i$ , indicate a path  $\pi \in \text{Paths}(T)$  for which  $\pi \not\models \varphi_i$ .

- Redo the previous task but ignore the fairness assumption.

### Exercise 4

(1 points)

We consider the release operator  $\mathsf{R}$  which is defined by

$$\varphi \mathsf{R} \psi \stackrel{\text{def}}{=} \neg(\neg \varphi \mathsf{U} \neg \psi) .$$

Prove the expansion law

$$\varphi_1 \mathsf{R} \varphi_2 \equiv \varphi_2 \wedge (\varphi_1 \vee \bigcirc (\varphi_1 \mathsf{R} \varphi_2)) .$$