

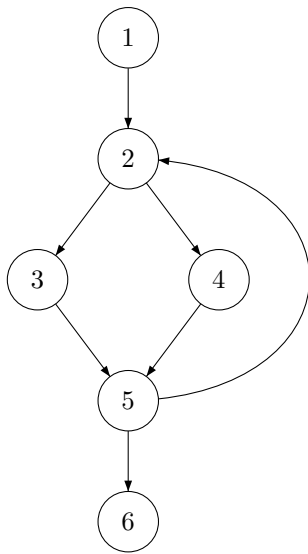
## 2. Exercise sheet *Static Program Analysis 2011*

Due Mon, 09. May 2011, *before* the exercise course begins.

### Exercise 2.1:

(3 points)

Both, *available expression* and *live variable analysis* had transfer functions of the form  $\varphi_l(d) = (d \setminus \text{kill}(B^l)) \cup \text{gen}(B^l)$ . Consider the following flowgraph and transfer functions on domain  $2^{\{a,b,c,d\}}$  and determine the minimal fixpoint using the worklist algorithm from the lecture.



$$\varphi_1(M) = (M \setminus \{a\}) \cup \{c\}$$

$$\varphi_2(M) = \begin{cases} M \cup \{b\} & \text{if } b \in M \vee c \in M \\ M & \text{otherwise} \end{cases}$$

$$\varphi_3(M) = \begin{cases} M \cup \{a, c\} & \text{if } (a \in M \vee d \in M) \wedge (b \in M \vee c \in M) \\ M \cup \{a\} & \text{if } a \in M \vee d \in M \\ M \cup \{c\} & \text{if } b \in M \vee c \in M \\ M & \text{otherwise} \end{cases}$$

$$\varphi_4(M) = \begin{cases} M \cup \{b, d\} & \text{if } b \in M \vee c \in M \vee d \in M \\ M \cup \{b\} & \text{otherwise} \end{cases}$$

$$\varphi_5(M) = (M \setminus \{a\}) \cup \{d\}$$

$$\varphi_6(M) = M$$

**Hint:** LUB = set union, relation = set-inclusion, initial value =  $\emptyset$ .

### Exercise 2.2:

(3 points)

Perform an available expression analysis on the following program using the meet over all paths (MOP) solution.

```

x := x + 1;
y := 1;
if x * y > y then
    x := x + 1;
    y := x + y;
else
    y := y + 1;
x := x + y;

```

### Exercise 2.3:

(2+2 points)

Consider domains of the form  $D = 2^M$  where  $M$  is a finite set. Transfer function  $\varphi_l : D \mapsto D$  is called distributive, iff for any  $d_1, d_2 \in D$ :

$$\varphi_l(d_1 \sqcup d_2) = \varphi_l(d_1) \sqcup \varphi_l(d_2)$$

- Show that any transfer function of the form  $\varphi_l(d) = (d \setminus \text{kill}(B^l)) \cup \text{gen}(B^l)$  is distributive in this setting.
- Show that the transfer function from Exercise 2.1 is not representable in *kill-gen* style, but still is distributive.

**Exercise 2.4:****(2+1+1 points)**

- (a) Perform a constant propagation analysis on the following program using the fixpoint iteration which was presented in the lecture (do not use the worklist algorithm!).

```
 $x := 10;$   
 $y := 1;$   
while  $x > 1$  do  
     $y := x * y;$   
     $x := x - 1;$ 
```

- (b) The aim of constant propagation analysis is to shift time needed for calculation from execution to compile time. For the given program with fixed variables  $x$  and  $y$  it would be possible to perform ALL calculations during compile time. Though constant propagation won't provide this result. Why?
- (c) Which of the program statements is the reason that constant propagation analysis is not distributive? Why? Give suitable variable valuations that confirm your explanation.