Semantics and Verification of Software

Lecture 3: Operational Semantics of WHILE II (Execution of Statements)

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Summer Semester 2010

Outline

1. Repetition: Evaluation Relations
2. Execution of Statements
3. Determinism of Evaluation/Execution

Evaluation of Arithmetic Expressions

Remember: $a ::= z \mid x \mid a_1\texttt{+}a_2 \mid a_1\texttt{-}a_2 \mid a_1\texttt{*}a_2 \in AExp$

Definition (Evaluation relation for arithmetic expressions)

If $a \in AExp$ and $\sigma \in \Sigma$, then $\langle a, \sigma \rangle$ is called a configuration.

Expression a evaluates to $z \in \mathbb{Z}$ in state σ (notation: $\langle a, \sigma \rangle \to z$) if this relationship is derivable by means of the following rules:

Axioms:

$$\frac{}{\langle z, \sigma \rangle \to z} \qquad \frac{}{\langle x, \sigma \rangle \to \sigma(x)}$$

Rules:

$$\frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1\texttt{+}a_2, \sigma \rangle \to z} \quad \text{where } z := z_1 + z_2$$

$$\frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1\texttt{-}a_2, \sigma \rangle \to z} \quad \text{where } z := z_1 - z_2$$

$$\frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1\texttt{*}a_2, \sigma \rangle \to z} \quad \text{where } z := z_1 * z_2$$

Evaluation of Boolean Expressions

Remember: $b ::= t \mid a_1 \texttt{=} a_2 \mid a_1 \texttt{>} a_2 \mid \neg b \mid b_1 \wedge b_2 \mid b_1 \vee b_2 \in BExp$

Definition (Evaluation relation for Boolean expressions)

For $b \in BExp$, $\sigma \in \Sigma$, and $t \in \mathbb{B}$, the evaluation relation $\langle b, \sigma \rangle \to t$ is defined by the following rules:

$$\frac{}{\langle t, \sigma \rangle \to t}$$

$$\frac{\langle a_1, \sigma \rangle \to z \quad \langle a_2, \sigma \rangle \to z}{\langle a_1 \texttt{=} a_2, \sigma \rangle \to \mathsf{true}} \qquad \frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1 \texttt{=} a_2, \sigma \rangle \to \mathsf{false}} \text{ if } z_1 \neq z_2$$

$$\frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1 \texttt{>} a_2, \sigma \rangle \to \mathsf{true}} \text{ if } z_1 > z_2 \qquad \frac{\langle a_1, \sigma \rangle \to z_1 \quad \langle a_2, \sigma \rangle \to z_2}{\langle a_1 \texttt{>} a_2, \sigma \rangle \to \mathsf{false}} \text{ if } z_1 \leq z_2$$

$$\frac{\langle b, \sigma \rangle \to \mathsf{false}}{\langle \neg b, \sigma \rangle \to \mathsf{true}} \qquad \frac{\langle b, \sigma \rangle \to \mathsf{true}}{\langle \neg b, \sigma \rangle \to \mathsf{false}}$$

$$\frac{\langle b_1, \sigma \rangle \to \mathsf{true} \quad \langle b_2, \sigma \rangle \to \mathsf{true}}{\langle b_1 \wedge b_2, \sigma \rangle \to \mathsf{true}} \qquad \frac{\langle b_1, \sigma \rangle \to \mathsf{true} \quad \langle b_2, \sigma \rangle \to \mathsf{false}}{\langle b_1 \wedge b_2, \sigma \rangle \to \mathsf{false}}$$

$$\frac{\langle b_1, \sigma \rangle \to \mathsf{false} \quad \langle b_2, \sigma \rangle \to \mathsf{true}}{\langle b_1 \wedge b_2, \sigma \rangle \to \mathsf{false}} \qquad \frac{\langle b_1, \sigma \rangle \to \mathsf{false} \quad \langle b_2, \sigma \rangle \to \mathsf{false}}{\langle b_1 \wedge b_2, \sigma \rangle \to \mathsf{false}}$$

($\vee$ analogously)

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Meaning of Statements

Effect of statement = transformation of program state

Example:

$$\langle \texttt{x := 2+3}, \sigma\rangle \rightarrow \sigma[\texttt{x} \mapsto 5]$$

where for every $\sigma \in \Sigma$, $x, y \in \mathit{Var}$, and $z \in \mathbb{Z}$:

$$\sigma[x \mapsto z](y) := \begin{cases} z & \text{if } y = x \\ \sigma(y) & \text{otherwise} \end{cases}$$

Execution of Statements

Remember:

$c ::= \texttt{skip} \mid x \texttt{ := } a \mid c_1 \texttt{;} c_2 \mid \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2 \mid \texttt{while } b \texttt{ do } c \in \mathit{Cmd}$

Definition 3.1 (Execution relation for statements)

For $c \in \mathit{Cmd}$ and $\sigma, \sigma' \in \Sigma$, the execution relation $\langle c, \sigma \rangle \to \sigma'$ is defined by the following rules:

$$\text{(skip)}\ \frac{}{\langle \texttt{skip}, \sigma \rangle \to \sigma} \qquad \text{(asgn)}\ \frac{\langle a, \sigma \rangle \to z}{\langle x \texttt{ := } a, \sigma \rangle \to \sigma[x \mapsto z]}$$

$$\text{(seq)}\ \frac{\langle c_1, \sigma \rangle \to \sigma' \quad \langle c_2, \sigma' \rangle \to \sigma''}{\langle c_1 \texttt{;} c_2, \sigma \rangle \to \sigma''} \qquad \text{(if-t)}\ \frac{\langle b, \sigma \rangle \to \mathsf{true} \quad \langle c_1, \sigma \rangle \to \sigma'}{\langle \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2, \sigma \rangle \to \sigma'}$$

$$\text{(if-f)}\ \frac{\langle b, \sigma \rangle \to \mathsf{false} \quad \langle c_2, \sigma \rangle \to \sigma'}{\langle \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2, \sigma \rangle \to \sigma'} \qquad \text{(wh-f)}\ \frac{\langle b, \sigma \rangle \to \mathsf{false}}{\langle \texttt{while } b \texttt{ do } c, \sigma \rangle \to \sigma}$$

$$\text{(wh-t)}\ \frac{\langle b, \sigma \rangle \to \mathsf{true} \quad \langle c, \sigma \rangle \to \sigma' \quad \langle \texttt{while } b \texttt{ do } c, \sigma' \rangle \to \sigma''}{\langle \texttt{while } b \texttt{ do } c, \sigma \rangle \to \sigma''}$$

An Execution Example

Example 3.2

- $c := \texttt{y := 1; while}\ \underbrace{\neg\texttt{(x=1)}}_{b}\ \texttt{do}\ \underbrace{\underbrace{\texttt{y := y*x}}_{c_1}\texttt{;}\underbrace{\texttt{x := x-1}}_{c_2}}_{c_0}$
- Claim: $\langle c, \sigma \rangle \rightarrow \sigma_{1,6}$ for every $\sigma \in \Sigma$ with $\sigma(\texttt{x}) = 3$
- Notation: $\sigma_{i,j}$ means $\sigma(\texttt{x}) = i$, $\sigma(\texttt{y}) = j$
- Derivation tree: on the board

Non-Terminating Statements

Corollary 3.3

The execution relation for statements is not total, i.e., there exist $c \in \mathit{Cmd}$ *and* $\sigma \in \Sigma$ *such that* $\langle c, \sigma \rangle \to \sigma'$ *for no* $\sigma' \in \Sigma$.

Proof.

Counterexample: $c =$ `while true do skip`
(by contradiction; see 1st Exercise) □

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Determinism of Execution Relation I

This operational semantics is well defined in the following sense:

Theorem 3.4

The execution relation for statements is deterministic, i.e., whenever $c \in \mathit{Cmd}$ *and* $\sigma, \sigma', \sigma'' \in \Sigma$ *such that* $\langle c, \sigma \rangle \to \sigma'$ *and* $\langle c, \sigma \rangle \to \sigma''$*, then* $\sigma' = \sigma''$*.*

The proof is based on the corresponding result for expressions.

Determinism of Evaluation Relations

Lemma 3.5

1. *For every* $a \in AExp$, $\sigma \in \Sigma$, *and* $z, z' \in \mathbb{Z}$*:* $\langle a, \sigma \rangle \to z$ *and* $\langle a, \sigma \rangle \to z'$ *implies* $z = z'$*.*
2. *For every* $b \in BExp$, $\sigma \in \Sigma$, *and* $t, t' \in \mathbb{B}$*:* $\langle b, \sigma \rangle \to t$ *and* $\langle b, \sigma \rangle \to t'$ *implies* $t = t'$*.*

Remarks:

- Lemma 3.5 is not implied by Lemma 2.6 (“$\sigma|_{FV(a)} = \sigma'|_{FV(a)} \implies (\langle a, \sigma \rangle \to z \iff \langle a, \sigma' \rangle \to z)$”)!

  The latter just implies
  $$\{z \in \mathbb{Z} \mid \langle a, \sigma \rangle \to z\} = \{z \in \mathbb{Z} \mid \langle a, \sigma' \rangle \to z\}$$
  while Lemma 3.5 states that
  $$|\{z \in \mathbb{Z} \mid \langle a, \sigma \rangle \to z\}| \leq 1.$$
- Lemma 3.5 can be shown by induction on the structure of expressions.

Excursus: Proof by Structural Induction IV

Application: Boolean expressions (Def. 1.2)

Definition: $BExp$ is the least set which

- contains the truth values $t \in \mathbb{B}$ and, for every $a_1, a_2 \in AExp$, $a_1 \texttt{=} a_2$ and $a_1 \texttt{>} a_2$, and
- contains $\neg b_1$, $b_1 \wedge b_2$ and $b_1 \vee b_2$ whenever $b_1, b_2 \in BExp$

Induction base: $P(t)$, $P(a_1 \texttt{=} a_2)$ and $P(a_1 \texttt{>} a_2)$ holds
(for every $t \in \mathbb{B}$, $a_1, a_2 \in AExp$)

Induction hypothesis: $P(b_1)$ and $P(b_2)$ holds

Induction step: $P(\neg b_1)$, $P(b_1 \wedge b_2)$ and $P(b_1 \vee b_2)$ holds

Proof (Lemma 3.5).

1. by structural induction on a (omitted)
2. by structural induction on b (omitted)

☐

Determinism of Execution Relation II

- How to prove that $\langle c, \sigma \rangle \to \sigma'$ is deterministic (Theorem 3.4)?
- Idea: use induction on the syntactic structure of c

Excursus: Proof by Structural Induction V

Application: syntax of WHILE statements (Def. 1.2)

Definition: Cmd is the least set which

- contains `skip` and, for every $x \in \mathit{Var}$ and $a \in \mathit{AExp}$, x `:=` a, and
- contains c_1`;`c_2, `if` b `then` c_1 `else` c_2 and `while` b `do` c_1 whenever $b \in \mathit{BExp}$ and $c_1, c_2 \in \mathit{Cmd}$

Induction base: $P($`skip`$)$ and $P(x$ `:=` $a)$ holds
(for every $x \in \mathit{Var}$ and $a \in \mathit{AExp}$)

Induction hypothesis: $P(c_1)$ and $P(c_2)$ holds

Induction step: $P(c_1$`;`$c_2)$, $P($`if` b `then` c_1 `else` $c_2)$ and
$P($`while` b `do` $c_1)$ holds

Determinism of Execution Relation III

- But: proof of Theorem 3.4 fails!
- Problematic case:

$$c = \texttt{while}\ b\ \texttt{do}\ c_0 \text{ where } \langle b, \sigma \rangle \rightarrow \mathsf{true}$$

- Here $\langle c, \sigma \rangle \rightarrow \sigma'$ and $\langle c, \sigma \rangle \rightarrow \sigma''$ require $\sigma_1, \sigma_2 \in \Sigma$ such that

$$\text{(wh-t)}\,\frac{\langle b, \sigma \rangle \rightarrow \mathsf{true} \quad \langle c_0, \sigma \rangle \rightarrow \sigma_1 \quad \langle c, \sigma_1 \rangle \rightarrow \sigma'}{\langle c, \sigma \rangle \rightarrow \sigma'}$$

  and

$$\text{(wh-t)}\,\frac{\langle b, \sigma \rangle \rightarrow \mathsf{true} \quad \langle c_0, \sigma \rangle \rightarrow \sigma_2 \quad \langle c, \sigma_2 \rangle \rightarrow \sigma''}{\langle c, \sigma \rangle \rightarrow \sigma''}$$

- c_0 proper substatement of c
  $\Longrightarrow$ induction hypothesis yields $\sigma_1 = \sigma_2$
- c not proper substatement of c $\Longrightarrow$ conclusion $\sigma' = \sigma''$ invalid!

Excursus: Proof by Structural Induction VI

Application: derivation trees of execution relation (Def. 3.1)

(skip): for every $\sigma \in \Sigma$, $\dfrac{}{\langle \texttt{skip}, \sigma\rangle \to \sigma}$ is a derivation tree for $\langle \texttt{skip}, \sigma\rangle \to \sigma$

(asgn): if s is a derivation tree for $\langle a, \sigma\rangle \to z$ (Def. 2.2), then $\dfrac{s}{\langle x := a, \sigma\rangle \to \sigma[x \mapsto z]}$ is a derivation tree for $\langle x := a, \sigma\rangle \to \sigma[x \mapsto z]$

(seq): if s_1 and s_2 are derivation trees for $\langle c_1, \sigma\rangle \to \sigma'$ and, respectively, $\langle c_2, \sigma'\rangle \to \sigma''$, then $\dfrac{s_1 \;\; s_2}{\langle c_1 ; c_2, \sigma\rangle \to \sigma''}$ is a derivation tree for $\langle c_1 ; c_2, \sigma\rangle \to \sigma''$

(if-t): if s_1 and s_2 are derivation trees for $\langle b, \sigma\rangle \to \textsf{true}$ (Def. 2.7) and, respectively, $\langle c_1, \sigma\rangle \to \sigma'$, then $\dfrac{s_1 \;\; s_2}{\langle \texttt{if}\ b\ \texttt{then}\ c_1\ \texttt{else}\ c_2, \sigma\rangle \to \sigma'}$ is a derivation tree for $\langle \texttt{if}\ b\ \texttt{then}\ c_1\ \texttt{else}\ c_2, \sigma\rangle \to \sigma'$

(if-f): analogously

(wh-t): if s_1, s_2 and s_3 are derivation trees for $\langle b, \sigma\rangle \to \textsf{true}$ (Def. 2.7), $\langle c, \sigma\rangle \to \sigma'$ and $\langle \texttt{while}\ b\ \texttt{do}\ c, \sigma'\rangle \to \sigma''$, respectively, then $\dfrac{s_1 \;\; s_2 \;\; s_3}{\langle \texttt{while}\ b\ \texttt{do}\ c, \sigma\rangle \to \sigma''}$ is a derivation tree for $\langle \texttt{while}\ b\ \texttt{do}\ c, \sigma\rangle \to \sigma''$

(wh-f): if s is a derivation tree for $\langle b, \sigma\rangle \to \textsf{false}$ (Def. 2.7), then $\dfrac{s}{\langle \texttt{while}\ b\ \texttt{do}\ c, \sigma\rangle \to \sigma}$ is a derivation tree for $\langle \texttt{while}\ b\ \texttt{do}\ c, \sigma\rangle \to \sigma$

Excursus: Proof by Structural Induction VII

Application: derivation trees of execution relation (continued)

Induction base: $P\left(\frac{}{\langle \texttt{skip}, \sigma\rangle \to \sigma}\right)$ holds for every $\sigma \in \Sigma$, and $P(s)$ holds for every derivation tree s for an arithmetic or Boolean expression.

Induction hypothesis: $P(s_1)$, $P(s_2)$ und $P(s_3)$ holds.

Induction step: it also holds that

(asgn): $P\left(\frac{s_1}{\langle x \texttt{:=} a, \sigma\rangle \to \sigma[x \mapsto z]}\right)$

(seq): $P\left(\frac{s_1 \quad s_2}{\langle c_1 \texttt{;} c_2, \sigma\rangle \to \sigma''}\right)$

(if-t): $P\left(\frac{s_1 \quad s_2}{\langle \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2, \sigma\rangle \to \sigma'}\right)$

(if-f): analogously

(wh-t): $P\left(\frac{s_1 \quad s_2 \quad s_3}{\langle \texttt{while } b \texttt{ do } c, \sigma\rangle \to \sigma''}\right)$

(wh-f): $P\left(\frac{s_1}{\langle \texttt{while } b \texttt{ do } c, \sigma\rangle \to \sigma}\right)$

Determinism of Execution Relation IV

Proof (Theorem 3.4).

To show:

$$\langle c, \sigma \rangle \to \sigma', \langle c, \sigma \rangle \to \sigma'' \implies \sigma' = \sigma''$$

(by structural induction on derivation trees; on the board) ☐