

**Exercise 1 (AM for While):**

**(1+1+1+1 Points)**

- a) Write a program in the WHILE programming language computing the greatest common divisor of two positive numbers given by variables  $x$  and  $y$ .
- b) Translate the program into intermediate code.
- c) Give a run of the program, i.e. the sequence of program states for  $x = 4, y = 2$ .
- d) Show that it actually does compute the greatest common divisor using formal semantics.

**Exercise 2 (AM: Repeat Until):**

**(1+1 Points)**

- a) Extend the WHILE language of the lecture with the construct **repeat**  $c$  **until**  $b$  and specify the corresponding translation function.
- b) When modeling the **repeat**  $c$  **until**  $b$  via the similar construct  $c; \text{while } \neg b \text{ do } c$ ; containing a while-loop in a straightforward way, the body  $c$  will be translated twice. Can you think of a way to avoid this double translation of  $c$ ?

**Exercise 3 (Compiler Correctness):**

**(5 Points)**

Provide the second/missing proof step for theorem of lecture 17, i.e. show that the following lemma holds:

**Lemma 1** For every  $c \in Cmd, \sigma, \sigma' \in \Sigma$  and  $e \in Stk$ ,  $\langle \mathfrak{T}_c[c], \epsilon, \sigma \rangle \triangleright^* \langle \epsilon, e, \sigma' \rangle$  implies  $\langle c, \sigma \rangle \rightarrow \sigma'$  and  $e = \epsilon$ .

You may use all theorems and lemmata presented in the lecture (except from the one to proof). Additionally you may find the following lemma useful.

**Lemma 2 (Decomposition Lemma)** If  $\langle c_1 : c_2, e, s \rangle \triangleright^k \langle \epsilon, e'', \sigma'' \rangle$ , then there exists a configuration  $\langle \epsilon, e', \sigma' \rangle$  and natural numbers  $k_1, k_2$  with  $k_1 + k_2 = k$  such that  $\langle c_1, e, \sigma \rangle \triangleright^{k_1} \langle \epsilon, e', \sigma' \rangle$  and  $\langle c_2, e', \sigma' \rangle \triangleright^{k_2} \langle \epsilon, e'', \sigma'' \rangle$ .