

Please hand in your solutions in groups of two or three students.
You can hand them in at the start of the exercise class, Monday April 29, 10:00 in AH 2.

Hints:

Single-step semantics for Boolean expressions: $\rightarrow_1^b \subseteq (BExp \times \Sigma) \times ((BExp \times \Sigma) \cup \mathbb{B})$

$$(Bt) \frac{}{\langle t, \sigma \rangle \rightarrow_1^b t}$$

$$(B = \text{true}) \frac{}{\langle z = z, \sigma \rangle \rightarrow_1^b \langle \text{true}, \sigma \rangle}$$

$$(B\neg) \frac{\langle b, \sigma \rangle \rightarrow_1^b \langle b', \sigma \rangle}{\langle \neg b, \sigma \rangle \rightarrow_1^b \langle \neg b', \sigma \rangle}$$

$$(B = \text{false}) \frac{}{\langle z_1 = z_2, \sigma \rangle \rightarrow_1^b \langle \text{false}, \sigma \rangle} (z_1 \neq z_2)$$

$$(B\neg\text{false}) \frac{}{\langle \neg \text{true}, \sigma \rangle \rightarrow_1^b \langle \text{false}, \sigma \rangle}$$

$$(B = 1) \frac{\langle a_1, \sigma \rangle \rightarrow_1^a \langle a'_1, \sigma \rangle}{\langle a_1 = a_2, \sigma \rangle \rightarrow_1^a \langle a'_1 = a_2, \sigma \rangle}$$

$$(B\neg\text{true}) \frac{}{\langle \neg \text{false}, \sigma \rangle \rightarrow_1^b \langle \text{true}, \sigma \rangle}$$

$$(B = 2) \frac{\langle a_2, \sigma \rangle \rightarrow_1^a \langle a'_2, \sigma \rangle}{\langle z_1 = a_2, \sigma \rangle \rightarrow_1^a \langle z_1 = a'_2, \sigma \rangle}$$

where $b \in BExp$; $z, z_1, z_2 \in \mathbb{Z}$; $a_1, a_2 \in AExp$. (For \geq, \vee, \wedge analogously.)

Exercise 1 (Single-Step Semantics):

(3 Points)

In the lecture we have defined a so-called *big-step semantics* for expressions, i.e., a relation $\rightarrow \subseteq (AExp \cup BExp \cup Cmd) \times \Sigma \times (\mathbb{Z} \cup \mathbb{B} \cup \Sigma)$ which yields the value of an expression within one step: $\langle (3 + 5) * (5 - 2), \sigma \rangle \rightarrow 24$. (Thus the intermediate results of the computation are “hidden” in the derivation tree.)

Alternatively it is possible to explicitly represent the intermediate steps by defining a *single-step semantics* $\rightarrow_1 = \rightarrow_1^a \cup \rightarrow_1^b \cup \rightarrow_1^c$, such that the following expression can be evaluated as:

$$\langle (3 + 5) * (5 - 2), \sigma \rangle \rightarrow_1 \langle 8 * (5 - 2), \sigma \rangle \rightarrow_1 \langle 8 * 3, \sigma \rangle \rightarrow_1 \langle 24, \sigma \rangle \rightarrow_1 24.$$

Give a complete specification of the single-step relation (see *Hints* above for an example of \rightarrow_1^b for Boolean expressions)

1. $\rightarrow_1^a \subseteq (AExp \times \Sigma) \times ((AExp \times \Sigma) \cup \mathbb{Z})$ for arithmetic expressions and
2. $\rightarrow_1^c \subseteq (Cmd \times \Sigma) \times (Cmd \times \Sigma)$ for commands.

Exercise 2 (Agreement of Big-Step and Single-Step Semantics):

(3 Points)

Show that the big-step relation and the single-step relation on arithmetic expressions, as defined in Exercise 1, are equivalent, i.e., that for every $a \in AExp$, $\delta \in \Sigma$, and $z \in \mathbb{Z}$:

$$\langle a, \delta \rangle \rightarrow z \quad \text{iff} \quad \langle a, \delta \rangle \rightarrow_1^* z.$$

Hint: use the existential quantifier \exists to assert that there exist z_1 and z_2 such that $z_1 + z_2 = z$.

Exercise 3 (repeat...until):**(1+2 Points)**

The **repeat...until** construct is similar to the **while...do** construct, except that the body of the loop is evaluated *at least once*. The loop ends when the guard evaluates to true. For example, with $\sigma(x) = 0$:

$$\begin{aligned} & \langle \text{repeat } x := x + 1 \text{ until } x = 2, \sigma \rangle \\ & \rightarrow \langle \text{repeat } x := x + 1 \text{ until } x = 2, \sigma[x \mapsto 1] \rangle \\ & \rightarrow \sigma[x \mapsto 2] \end{aligned}$$

- (a) Define the execution relation for the **repeat...until** construct without using the **while...do** construct.
- (b) Show that **repeat** c **until** $b \sim c; \text{while } \neg b \text{ do } c$ for any $b \in BExp, c \in Cmd$.

Exercise 4 (Infinite Loop):**(2+1 Points)**

- (a) Show that the following proposition holds: for all states σ, σ' , $\langle \text{while true do skip}, \sigma \rangle \not\rightarrow \sigma'$.
- (b) Show that in the single-step semantics you defined in Exercise 1, the execution of $\langle \text{while true do skip}, \sigma \rangle$ does not terminate.